Forensic Analysis of Volume Shadow Service ($RecycleBin) of Win10

by

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CERTIFICATE

This is to certify that **NS Nosheen Manzoor** Student of **MSIS-14** Course Reg.No **00000119618** has completed her MS Thesis title **"Forensic Analysis of Volume Shadow Service ($RecycleBin) of Win10 "**under my supervision. I have reviewed her final thesis copy and I am satisfied with her work.

Thesis Supervisor

(**Assistant Professor Mian Muhammad Waseem Iqbal**)

Dated: _____June 2019

Declaration

I hereby declare that no portion of work presented in this thesis has been submitted in support of another award or qualification either at this institution or elsewhere.

Dedication

“In the name of Allah, the most Beneficent, the most Merciful"

I dedicate this thesis to my Mother, Father, Husband and Teachers who supported me in each step of the way.

Acknowledgments

I am thankful to Allah the Almighty, for giving me the courage and strength to complete my thesis. I would also like to pay special thanks to my Supervisor, Assistant Professor Waseem Iqbal, for his continuous support and guidance throughout this thesis and believing in me. Also, I would thank my committee members; Lecturer Narmeen Shafqat, and Asst Prof Waleed Bin Shahid for their support and knowledge regarding this topic.

I am also thankful to my Family especially to my Father, Mother and my Husband for sparing me from family commitments, motivating me and praying for my success. Without their support, it would have not been possible for me to complete this research work. Last but not the least, my special gratitude to all my wonderful teachers of Department of Information Security at Military College of Signals, for imparting knowledge and enabling me to complete this task.

Abstract

In Todays' digital world everything is shifting into the smart technology and people have started relying on this digital world. Size of the storage media has also been increased day by day. When everything is shifted to the digital world therefore crime has also been shifted to the Digital crime. In digital crime finding out the digital evidence from the storage media is becoming complex and time consuming.

It is better and interesting for the investigator to start carving for the evidence from the most crucial areas like windows Volume Shadow Service.

Volume Shadow Copy is considered as gold mine for the forensic investigator as it generates differential backups. Previous versions of the files, Recycle bin and state of the $logfile get saved in Volume Shadow Copy which holds clumps of crucial data for the investigators.

Volume Shadow Copy lets the investigator to understand the state of the system on a particular date. Whatever is deleted from the system even its deleted permanently with a wipe utility may have its presence in VSC (Volume Shadow Copy).Being not accessible to the user in normal environment and being "Read only" in nature preserves the evidences to a great extent.VSC in series gives the idea of routine and activities performed by the accused in a sequential manner. An experiment with two case scenarios e.g Case scenario1, stolen financial information and Case scenario 2, modified health information of a patient has been conducted to prove the importance of the Volume Shadow Copy. Methodology has been proposed to extract the data from the Volume Shadow Copy of Windows 10 to find the evidence from Volume Shadow Copy's store which gives access to the previous version of not only the user files but also from the system files.

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Chapter1

Introduction

1.1 Overview

Volume Shadow Service is a Microsoft's built-in service of windows. It was first introduced in Windows server 2003 named as "Shadow Copies for Shared Folder". When this service was added in windows vista then its name was changed to "Volume Shadow Copies" in 2007.Volume Shadow copy creates differential backup copies of crucial data. Volume Shadow Service provides the facility to use previous version in order to find your data which may corrupt by any virus or accidental data loss.[1][2]

Volume shadow copy is a Microsoft window's inbuilt service that permits the user to take backups of the system either its manual or automatic even when the system is in use. The block size of the hard drive of the system is 16KB. These blocks of the computer system are constantly monitored and as soon as any modification or alteration took place in a block window start implementing this modification in the block of volume shadow service after storing it on the storage location. The window creates a backup of the block in this fashion. System's settings and all important data of drive C records in Volume Shadow Copy Service. It enables the system to encounter unpredictably data deletion and from events which destabilize the system, like a virus attack or the inaccurate installation of a software package or any other hardware device. It's been remarked as a gold mine of forensic proof because of the amount of data it records.[2]

This information is very valuable from the forensic point of view due to the following reasons.

a. It lets the investigator understand the state of the system on a particular date.

b. Whatever is deleted from the system even its deleted permanently with a wipe utility may have its presence in VSC (Volume Shadow Copy)

c. Being not accessible to the user in the normal environment and being “Read-only” in nature preserves the evidence to a great extent.

d. VSC in series gives the idea of routine and activities performed by the accused in a sequential manner.

There are so many built-in features available in Windows 10 that can be utilized as a good source of forensic artifacts and Volume shadow service is from one of them. Shadow copies are created in two different ways, a differential backup or complete backup. Complete backup generates a complete copy of data available on disk. In the differentialbackup, only those changes are backed up which are made in a specified block on the original volume. To keep track of the logical construction of the volume shadow copy a buffer is used. [3]

Automatic and manual creations are two different ways to create volume shadow copies. In automatic creation volume shadow copy is generated or activated by the operating system as soon as new software is installed and when an update of the system is installed. Volume Shadow service then creates an image of only those files which are changed since the last backup. Users create manual volume shadow copy when they make changes in their files and wants to trigger volume shadow copy manually. Creation of Volume Shadow service is a three-step process. Freeze: Hard disk of the computer marked as read-only Snap: Image of the system/Hard drive created.Unfreeze: Hard disk of the system get released and volume shadow service runs in the background.[4]

There are three main components of Volume Shadow Service Writer: Volume Shadow writer is responsible for informing the backup device that how to back up the information, applications and their data. Requestor: A Volume Shadow Service

requestor is responsible for starting the VSC processes. Mostly VSS requestors are

backup applications. Provider: VSS provider acts as a middle layer between backup processes, operating system,and hardware.[2]

1.2 Motivation and Problem Statement

Volume Shadow Copies provides extra information with extra data that normally not available on the system.It enables the forensic investigator to know about what was happening in the system before he/she may start the investigation. It provides the point in time copies of the user as well as system data. Shadow copy is the vigilant tool to recover the previously permanently deleted files by the user.

For every newly released operating system, Forensic investigator must re-consider that new version of the operating system to determine any minor/major changes which may affect their investigation. Forensic Analysis of Windows 7 has been done but information available in the literature is not sufficient to analyze the new versions of the operating systems like windows 10.

1.3 Objective

The main objectives of the thesis are:-

a. Forensic Analysis of $Recyclebin from volume shadow service of Microsoft windows10
b. Comparison between $Recyclebin from volume shadow service of Microsoft window7 and win10

1.4 Thesis Contribution

To best of my knowledge, limited research has been done on forensic analysis of Windows 10 and especially how to extract the files and folders from the difference files. Most of the means which I have explored are blogs, presentations and

articles have only given very basic knowledge about forensic analysis of Windows 10. Moreover, the internal structure of the shadow has not been discussed.

The main contribution of this research work are as follows

a) We have proposed a mechanism to identify the Volume Shadow Copy stores of Windows 10 from the system image.

b) We have discussed the internal structure of the volume Shadow Copy of Windows 10.

c) We have proposed the way to recover the files from the Volume Shadow Copy of Windows 10.

1.5 Thesis Organization

The thesis is structured as follows:

- Chapter 2 contains the literature reviewed in the thesis. The general introduction of the Volume shadow copy, working of Volume Shadow Copy, Creation process of Volume Shadow Copy, Access methods of Volume Shadow Copy used for Windows 7,Vista
- Chapter 3 contains the test beds, Experimental test case scenarios. proposed methodology for accessing the Volume Shadow Copy
- Chapter 4 Analysis of $Log file and $ Recycle.Bin, their importance and artifacts.
- Chapter 5 contains the Results and analysis reports, Winhex screenshots of experiments
- Chapter 6 contains the comparison of different freely available tools for the analysis of the Volume Shadow Copy.
- Chapter7 contains the discussion, conclusion of the thesis and Future work.

Chapter2

LiteratureReview- UnderstandingVolume Shadow Copy

In order to forensically examine volume shadow copy of windows 10 in a better way, it is important to explore research papers and articles. Unluckily limited research has been found on forensic analysis of Windows 10 and especially on forensic analysis of $Recycle bin of Volume shadow copy of Windows 10. The information available in the shape of research done by the research community. Most of the means which I have explored are blogs, presentations and articles have only given very basic knowledge about forensic analysis of volume shadow copy of windows 10.

2.1 Understanding Volume Shadow Copy

In order to understand how Volume Shadow Copy works it is important to understand its layout, structure, and configuration. This chapter discusses all necessary concepts Volume Shadow Copy, related to this research thesis.

2.1.1Volume Shadow Copy Volume Shadow Copy Service (VSS) is a Component Object Model (COM) interfaces in Microsoft Windows built-in service to perform volume backups. Excellent coordination is required between backup application, user application which is going to be backed up and hardware and software management. Volume shadow copy which was first introduced in 2003 provides the coordination between these applications.

Volume Shadow Copies are transparently maintained by the Windows. Volume Shadow Copy operates at a lower layer than NTFS as shown in fig1 [9].

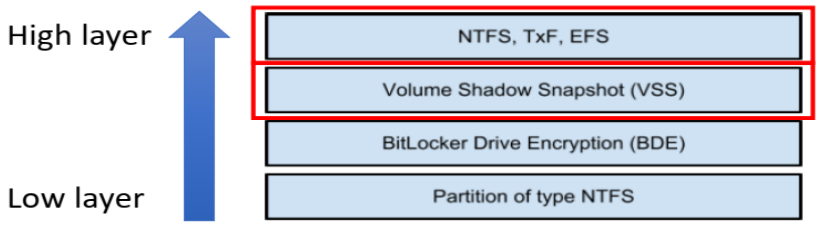


Fig:1 Layer of NTFS Volume

2.1.2 Working of Volume Shadow Copy: Volume Shadow Copy consists of thefollowing basic components.

VSS service: It coordinates with all other components of the VSS to create shadow copies smoothly

VSS Requestor: It is backup software that requests the operating system to perform a backup.Window server backup utility basically performs this duty.

VSS Writers: This software assures that the consistent backup copies of the windows are generated. They provide data integrity during backups process.

VSS Provider: VSS Providers can be VSS hardware and software and they actually generate shadow copies and then work on their maintenance as well.[2]

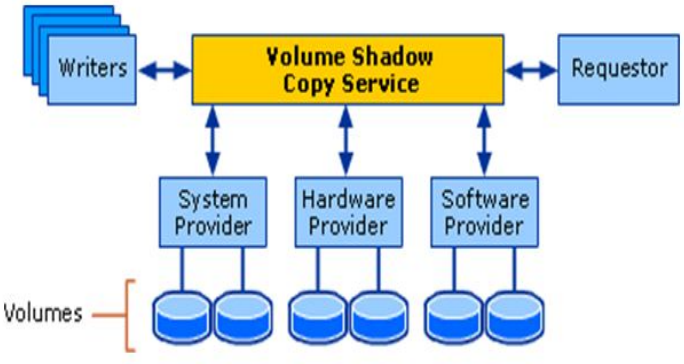


Fig 2: Architectural diagram of Volume Shadow Copy

2.1.3Creation Process of a Volume Shadow Copy

In the creationprocess, all the components of VSS service work together with a high level of coordination. Shadow copy creation process is shown in fig. 3.

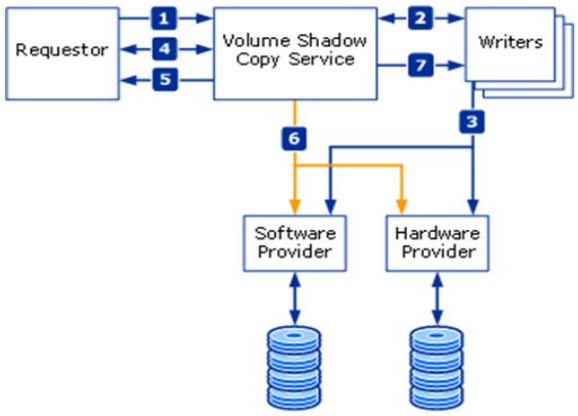


Fig3: Creation Process of Volume Shadow Copy

- Volume shadow copy service specify the writer after getting metadata for preparation of shadow copy created on the request of VSS Requestor
- XML description has been created by every writer involved in the creation of volume shadow copy
- Each writer creates an XML description of the components and data stores that need to be backed up and provides it to the Volume Shadow Copy Service. The writer also defines a restore method, which is used for all components. The Volume Shadow Copy Service
- provides the writer's description of the requester, which selects the components that will be backed up.
- All the writers then notified by the Volume Shadow service to be ready to create a shadow copy.
- All the applications are temporarily frozen for less than 60 seconds to write-up the data into the shadows.

- Within 10 seconds shadow has been created then all the I/O operation related to different applications gets released for their normal working.

2.2 Enabling Volume Shadow Copy

In order to enable the Volume Shadow Copy, we have to turn on the system Protection feature from the control panel.

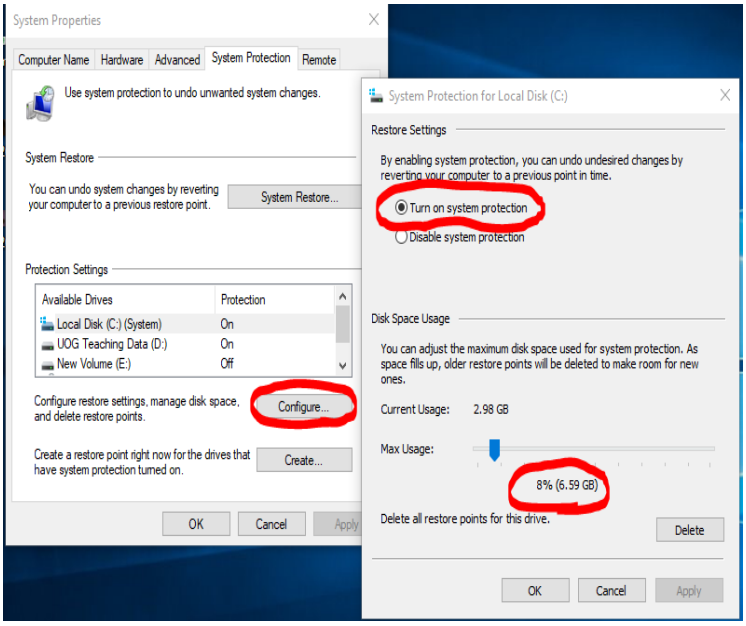


Fig4: Configuration of Volume Shadow Copy

Older shadow copies are deleted by the window when window run out of space. Volume Shadow Copy of a specific volume is stored in the volume itself so if the volume gets corrupted the Volume Shadow Copy of that volume will also be corrupted. Volume Shadow Copy is a block level incremental backup. Block size to be increment is 16 KB. Data of Volume Shadow Copy cannot be changed until or unless it is deleted. Volume Shadow Copy generates incremental backups and provides the facility to the user to restore the system's previous state when required. [5]

GUID{3808876b-c176-4e48-b7ae-04046e6cc752} is used by the Volume Shadow Copy to distinguish the attributes of the shadow copies which includes header files

and store files.

2.3 Method of Storing Data in Volume Shadow Copy

As soon as volume shadow copy is created its corresponding catalog and store is allocated to the volume shadow copy. Data stored in 16KB blocks. As an example, 60KB file will be stored in 4 data blocks as I data block is of 16 KB shown in fig5a.[6]

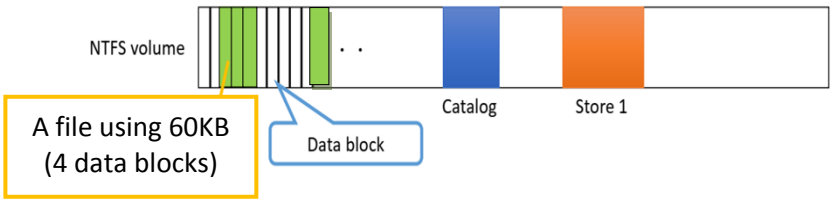


Fig: 5a Data storage method in volume shadow Copy

As soon as the data in the specific block is changed it's been copied to the store1 as shown in the fig 5b.[6]

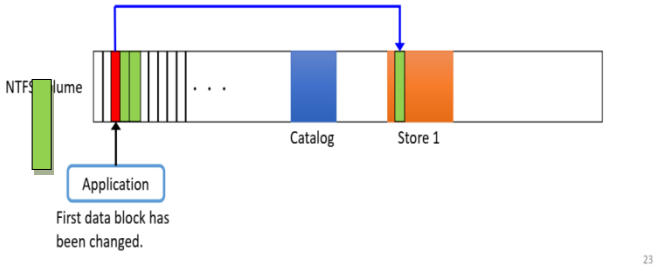


Fig: 5b Data storage method in volume shadow Copy:

When we create a second Volume Shadow Copy second store is allocated to the volume shadow copy and second entry has been made into the catalog entries. Now only those Blocks have been backed up which have been changed since last Shadow Copy [6].

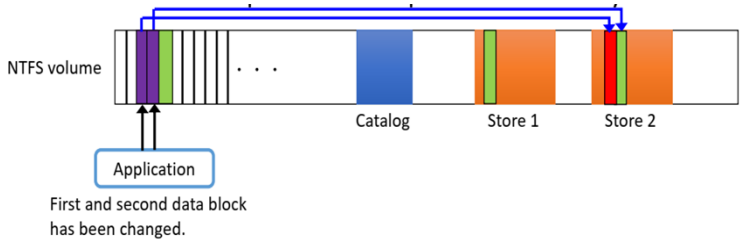


Fig5c:Data storage method in volume shadow Copy

2.4 Location and Structure of Volume Shadow Copy

Following is a layout through which windows Operating system access the volume snapshot by accessing the volume snapshot header.

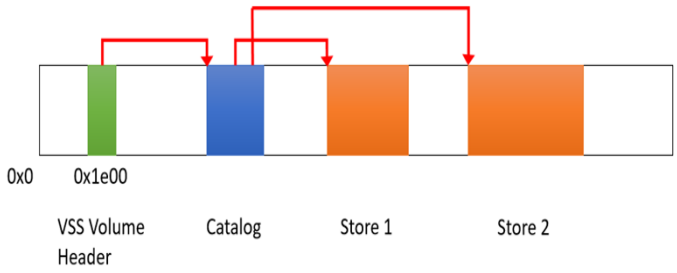


Fig 6: Volume Shadow Copy working layout

NTFS volume header contains the Volume Shadow Copy header as its part. Volume header data always starts at offset 7860(0x1e00) in windows vista, 7, 8 and in Windows 10 as well. Its size is 512 bytes equals to one sector [6].

2.4.1Structure of VSS Volume header

Volume Shadow Service is basically is located at offset 7680(1x1E00) of an NTFS volume. Its size is 512byte =1sector. Volume snapshot header consists of the Volume Shadow Copy identifier and offset of the first catalog blocks shown in table1 [7].

Table 1: Catalog Block Header[7]

0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
Volume Shadow Copy Identifier															
Version		Record Type(0x01)						Current offset relative to the start of the volume							
Unknown (Next offset relative to the start of the volume) 0x1e00								Unknown empty values							
Catalog offset relative to the start of the volume (value=0 if there is no catalog)								Maximum size (set to 0 if unbounded)							
Volume identifier (contains GUID)															
Shadow copy storage volume identifier (GUID)															
Unknown		Unknown empty values													

2.4.2Catalog Block Header

The catalog holds the information about each and every snapshot. Every catalog has One or more than one catalog blocks. Each catalog block consists of catalog block header and catalog block entry as shown in table 2. If Volume Shadow Copy is enabled but has no snapshot, then no catalog will exist [7][8].

Table2: Catalog Block Header[7][8]

0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
VOLUME SHADOW COPY Identifier (contains GUID)															
Version 0x01				**Record Type(0x02)**				Relative catalog block list offset (The offset is relative to the start of the catalog block)							
Current catalog block list offset relative to the start of the volume								Next Catalog block offset							

Table 3: Catalog bock header types

Version	Type
0x01	Windows Vista, and 7
0x02	Windows 10

2.4.3Catalog Entries

Catalog entry started directly after the catalog block header in this case offset of first catalog entry is (0x71f0000).Catalog entry type0x03 is found directly after the catalog entry type of 0x02[8]. If the catalog entry type 0x03 is present, then it shows that stores are present on the volume. If a system hasfour-volume shadow copies it means four 0x02 and 0x03 entry type(2 entries for each Volume Shadow Copy=8 entries)

2.4.3.1Structure of Entry type 0x02

Structure of entry type 0x02 is described in table 4 [7][8].

Table4: Structure of Entry type 0x02[7][8]

Entry type 0x02																
	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
	Catalog entry type								Volume Size							
	Store identifier (contains GUID used in store filename)															
	Unknown sequence number								Flag values (440 in windows 10) 40 in win 7,8							
	Shadow copy creation time								Unknown Empty values							

2.4.3.2 Structure of entry type of 0x03

Structure of entry type of 0x03 described in table 5.

Table5: Entry type 0x03[10]

Entry Type 0x03	
Store identifier (contains GUID used in store filename)	
Store header offset relative to the stat of the volume	Store block range list offset relative to the start of the volume
Current bitmap offset relative to the start of the volume	NTFS metadata file reference
Allocated size	Store previous bitmap offset

2.4.4Store (Actual volume snapshot)

The actual data blocks of Volume Shadow Copy are stored in stores. Following

data structure keeps track of the volume snapshot locations.

2.4.4.1Store Blocklist

It contains store block header of type 3 and is the size of 128 bytes followed by 32 bytes of index shown in table 6

Table 6: Store blocklist [7][8]

0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
Volume Shadow Copy Identifier contains GUID															
Version 0x01				Record Type (0x03) Store Header				Relative block offset							
Current block offset								Next block offset (0 incase of the last block)							
Size of store information								Unknown empty values							

2.4.4.2Catalog Block List

It starts directly after the store. It consists of original data block offset and relative store data block offset and store block descriptors.

2.4.4.3Store block descriptor

Flag 0x00	**Normal descriptor**
Flag 0x88	Complete block of the snapshot will be mapped to the original block
Flag 0x01	**Overlay descriptor** In an overlay a descriptor,, bitmap allocation table contains the data about to fill the block.
Flag 0x02	**Forward descriptor** The relative offset will be mapped to the original offset to the next block
Flag 0x04	**Invalid index record** The block will be ignored if the flag set to 0x04

It consists of store block header of type3 and stores block descriptor of the size of 32 bytes. Store block descriptor consists of original block offset as well as relative store data block offset. Original data block offset should be replaced by the shadow copy data block. The mapping between snapshot data and original is based on flag fields. Values of the flags can be interpreted as follows [7][8].

2.5 Accessing Volume Shadow Copy

Volume Shadow Copy creates differential backups as it only stores the changed part of the file, not a complete file. In order to access files from the Volume Shadow

Copy, all the stores and 16 KB data blocks will be read and combined with the live volume to recreate the complete file [6] [7][9].

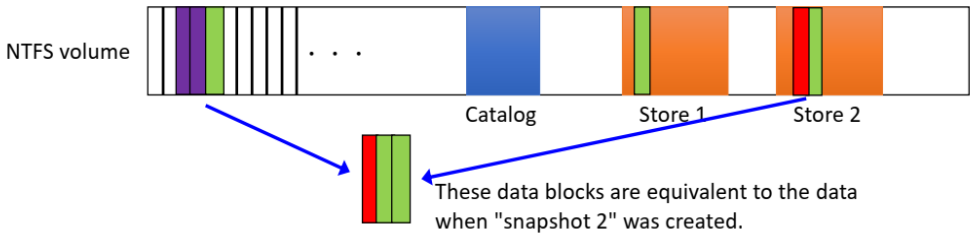


Fig:7 Reconstructing Volume Shadow Copy: These data blocks are equivalent to the data when second shadow was created

Reading of the data has been started from the most recent store to the oldest one. As in Example described in fig. data block from the current volume is combined with the data block of store reproduces the file which is equal to the file when second shadow copy has been created [7][6][8].

As we have created only two shadow copies, we have only two stores. The file reconstructed from the previous step further will be combined with the data block of stored in the first step reproduced the file equivalent to the file when shadow copy 1 was created.as shown in fig 7a

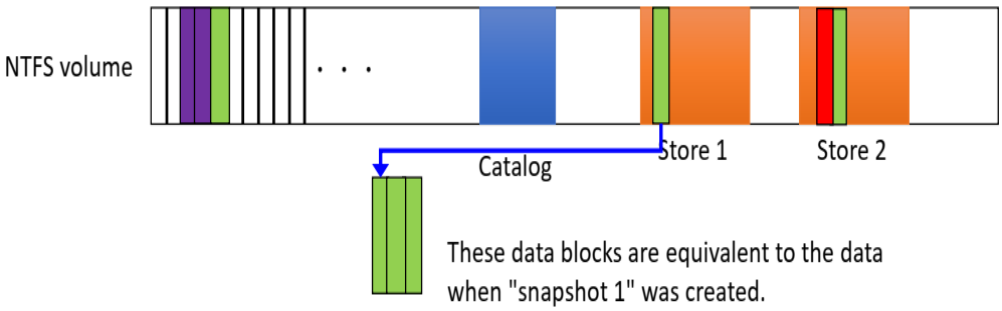


Fig7a: Reconstructing Volume Shadow Copy: These blocks represent the data when 1st shadow copy was created

Chapter3

Test Beds

This Chapter provides the details of testbed and experiments conducted for this research. Testbeds, Test Cases, image files and proposed forensic methodology has been discussed in this chapter

3.1 Detail of Test image

A test image has been created for experimental purposes. Windows 10 Operating system was used. Details are as under:

Victim Machine

The experiment was conducted on Windows 10, Intel i5 processor, 4GB Ram, and 520 GB hard disk.

Forensic machine

The investigation was conducted on Windows 10, Intel i7 processor, 8GB Ram and 2TB hard disk.

3.2 Proposed Forensic Methodology

The main purpose of this research is to find artifacts from volume shadow copy, especially from $Recycle.Bin and $logfile to prove the importance of the Volume Shadow Copy for the forensic investigators. For a complete analysis of the Volume Shadow Copy, two different case scenarios have been generated and analyzed with different volume shadow copies. Different files were created modified and deleted in order to generate evidence. After analyzing the volume shadow copy, results will be compared with existing work.

The experiment consists of two case scenarios and the experiment consists of the

following steps: 1) Creation2) Extraction3) Analysis of the evidence and then comparison with the previously done work.

1st Case Scenario: In the first case scenario financial information of a client of a company has been stolen by their own employee. Later on, during investigation, it has been found that suspects have copied those stolen files on their personal computer. The forensic investigator has to dig out and find proof of stolen information. For this purpose, we will create one file from scratch and two files have been downloaded from the internet.

A file named as Account_info.txt has been created from scratch with information of the account holder.2nd File check_bank.jpg has been downloaded from the internet and placed it on the desktop of the suspect's computer. The 3rd file is Bank_Statement.jpg initially has been placed on desktop and shadow copy has been created, then these files were moved from desktop to the documents in Account_info folder. At last step Folder named as Account_info has been deleted from the system and another shadow copy has been created. All these four shadow copies will be manually analyzed to prove the importance of the shadow copies for the forensic investigators.

2nd Case Scenario: Health information of a patient has been stolen, now the investigator has to find out the evidence from suspect's computer. For this situation, three files Blood_Count.docx, Brain_Scan.jpg and Case history.pdf will be copied on desktop of the suspects computer. Blood_Count.docx has been modified and changes have been saved to the original file. Brain_scan has also been modified and saved as Brain_Scan.png. after making changes and using the files these files will be deleted from the suspects computer.

A forensic image of the suspect's computer has been taken with the help of FTK imager named as **Financial_001**. Four Shadow Copies has been created for the

Case1:

Volume Shadow Copy I: Contains the shadow copy of the whole system after copying the stolen information into the suspect's computer. Initially, files were saved on desktop.

Volume Shadow Copy II: The second shadow was triggered when the files were moved into the new folder account-info created on the desktop.

Volume Shadow Copy III: the third shadow was created when the folder Account_info was moved to the Document folder

Volume Shadow Copy IV: Fourth shadow copy has been creating when the folder Account_info has been deleted. In order to traces, the changes shadow copy has been composed using the steps discussed in section 3.3

3.3 Experimental Testbeds

A forensic image of the system has been taken with the help of FTK Imager. After successfully taken the image it was analyzed by WinHex18 academia version and for verification purposes, automated analysis with the help of different freely available forensic investigation tools has been used which gives the support to the volume shadow copy. To conduct the experiment following tools and technologies have been used.

3.3.1 FTK Imager

FTK (Forensic Tool Kit) imager is a tool to create a forensic image of the disk as a whole or in parts that may be reconstructed at the end by Access Data. It also creates MD5 Hash values of the image for verification purposes. [18]

3.3.2 WinHex

In this research WinHex, 15.2 Academia version has been used. Winhex is a

universal hex editor which is helpful in computer forensic investigation, low-level data extraction. It gives the facility to carve the data. Extract and analyze all kinds of files. It gives the facility of refining volume shadows which automatically mount the Volume Shadow Copy to help the investigator. Refining volume shadow feature is not included in academia version, so we investigate our case without using this feature. In this investigation academia version has been used and it does not include refine volume shadow feature [19].

3.4 Proposed methodology of reading data from Volume Shadow Copy of Windows 10

Complete workflow of the research has been shown in fig 8.

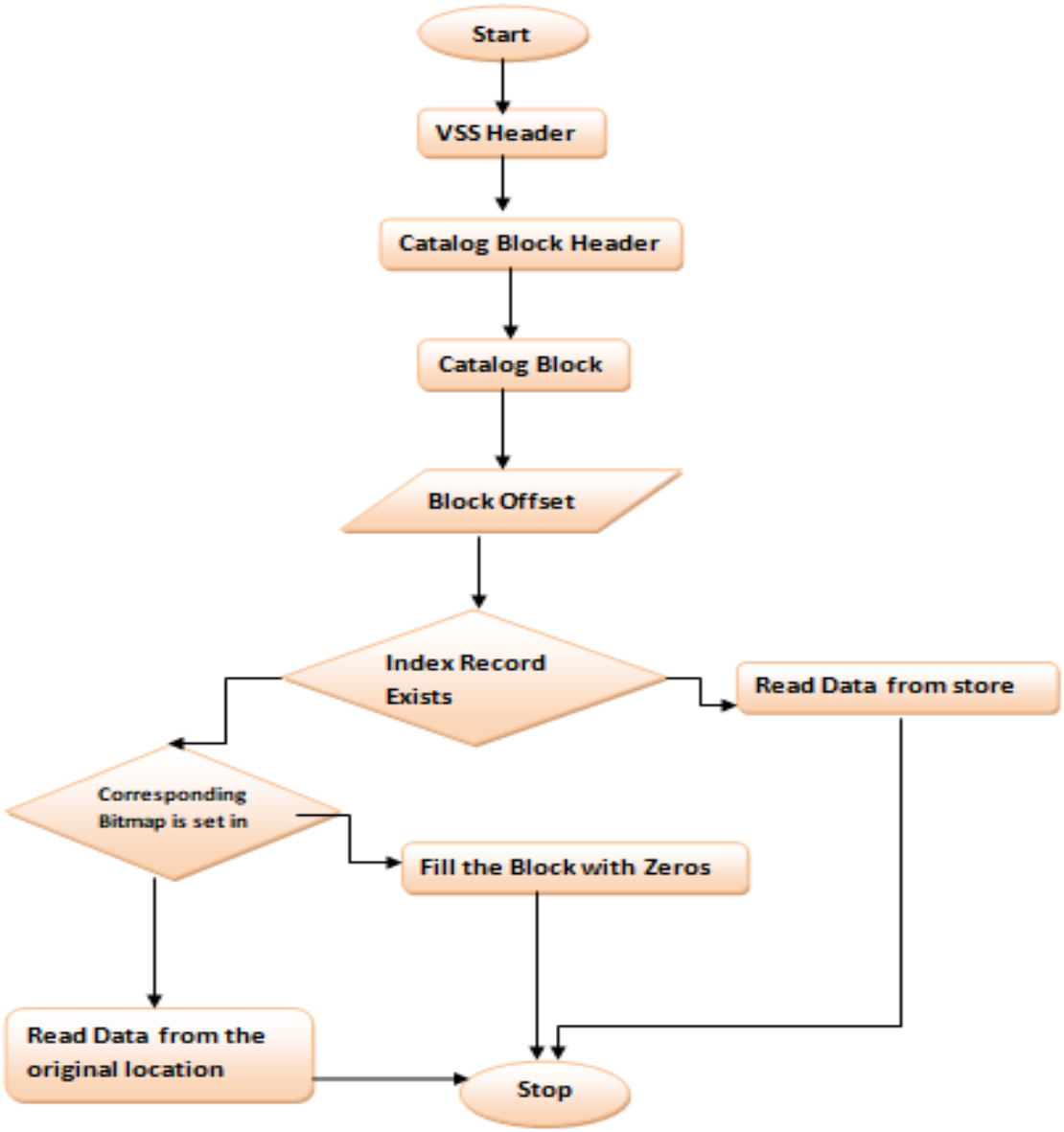


Fig8: Flowchart of reading data from Volume Shadow Copy

Chapter 4

$LogFile and $Recycle.Bin Analysis

There is a possibility that no artifacts of a deleted file exist in $MFT as it can be overwritten by new files. As all Volume Shadow Copies starts with $logfile which provides crucial information about the evidence. This research paper [10] also focuses on $logfile and discuss it in detail. $logFile has been discussed in detail which includes all types of records, the structure of the records and all information which is logged in them.

From the first 16 entries, $logfile is situated in a 2ndposition [11]. $logFile is a value-based log, recording changes of the NTFS file system. Default page size for Records is 4096 bytes or 0x1000 (which can be increased or decreased). Each record has a unique LSN ($logfile Sequence Number) which increments each time the file has been used. Log File Sequence Number is used to correlate file record of $logfile and $MFT. Each $MFT file contains LSN from 0x08-0x0F [11].

Two transactions are used by NTFS to complete filing tasks: In first transaction, files have been updated and in second transaction attributes of the $MFT has been updated. If the system fails after the first transaction, then the $Log file is used to recover the system. Every transaction is recorded in a $Log file. Every activity, e.g. renaming of a file, deleting a file have many transactions linked to complete the transaction. Therefore, a chain of operational records has been generated in $logfile to complete that activity. Because of this reason $logfile is important for forensics. Operation records of $logfile sustain data before transaction (for restoration/rollback/Undo) and data after transaction (Redo). Incase of Renaming a file following type of information is maintained [12].

4.1 Types of Log Records.

There are two types of records are recorded in two primary zones; Restart Area Records and Logging Area Records [13].

4.1.1 **Restart Area Records**. In this area, two records have been stored and both starting with "RSTR" each of length 0x1000.The 2nd record is the copy of the first record. Current LSN record in this record holds the information of last operation record. Structure of Restart area record is shown in table [13].

4.1.2 **Logging Area Records**. Real operational records are stored in the Logging are [13]. Normal page and buffer page area are two main divisions of the Logging area record. Logging area is divided into buffer Page Area and Normal Page. Buffer Page Area (0x200 to 0x4000) consists of first two pages. The second page is the copy of the first page. Last operation record is stored in Buffer Page area. Older record is pushed into the normal page area when Buffer Page area has been full. Normal Page area extends from 0x4000 to end of $logfile shown in the table7.

Structure of Page. Every page in the Logging area consists of a header followed by more than one operational records. Page header includes data of that page and its structure is shown in table 8 [13] with an example shown in fig 9.

Table 7: Start of the Log record

0x00	"RSTR" After checkdisk.its "CHKD"		UpdateSeq.ArrayOffset	UpdateSeq.ArraySize
	Check DiskLSN (Multi Sectorheader,allzero less "RSTR"changeto "CHKD",lastLSNfoundbycheck disk)			
0x10	System PageSize (fatalerror if System PageSize≠Logpage)		Log Page Size	
	RestartArea Offset (from"RSTR")	Minor Version (-1= beta,0=Transition, 1=updateseq. sp.)	Major Version (-1= beta,0=Transition, 1=updateseq. sp.)	
0x20	UpdateSeq.Array			
0x30	CurrentLSN (currentlogicalendof the logfileto facilitaterestart)			
	LogClient (maxclient sp.for thislog file)	ClientFreeList	ClientIn-useList	Flags
0x40	Seq.Number bits		RestartAreaLength	ClientArrayOffset (fromthe startof thisstructure)
	FileSize of $Logfile			
0x50	LastLSNdatalength(excludingR estartpageheader)		RecordPage Header Length	Log PageDataOffset
	RestartOpenLogCount (logfileopencount, todeterminethe changein thedisk e.g. remounting)		Padding	
0x06	Log clientArray (ClientData)			
	(Cont.)Log clientArray (ClientData)			
0x07	OldestLSN (Required tobein the logfilebythisclient)			
	ClientRestartLSN (LSNoflatestclient restartareawritten to thedisk,generallyCurrentLSN))			
0x08	Previous client(0xFF meansnoclient)	Nextclient (0xFFmeans no client)	Seq.Number(incremented on recordre-use)	Align Word (alignment field)
	Alignworld(alignt heentire record)		ClientNamelength (always 8)	
0x90	Client name (NTFSwithremainingbytes setto zero)			

Table8: Page Header of operational Record

	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E
0x00	"RCRD" (signature)				Update Sequence Offset		Update Sequence Count		Last LSN orFile Offset (Last overallLSN on the page, including overlapping LSNs to next page)						
0x10	Flags				Page Count (No. of pagesused for transaction run)		Page Position (Current page no.)		Next Record Offset (Next LSN on the page)		Word Align		DWordAlign		
0x20	Last EndLSN (Last completedLSN on the page)														
0x30	Update SequenceArray (Arraycontaining theupdate seq. numberforreplacement First,two bytes of thevalue isthe UpdateSeqValue, usedevery512 bytes Followed by updatesequencearrays fora number of two-bytegroups definedinupdatedseq.countless 1)														

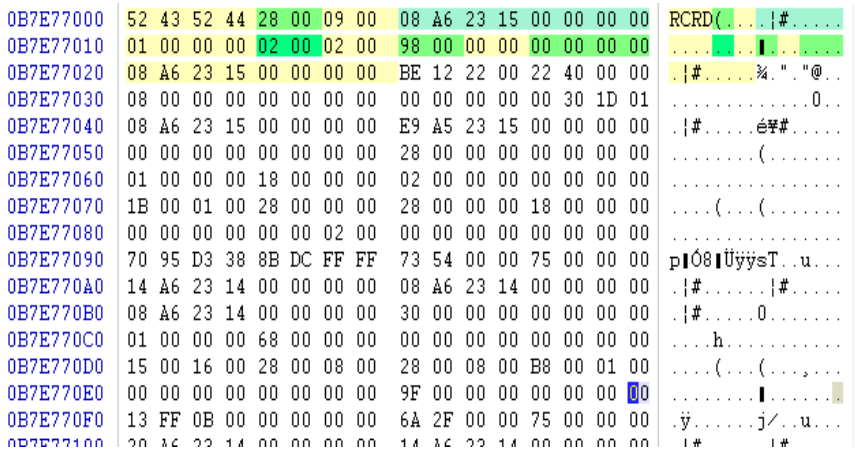


Figure 9:Example of Page Header of $LogFile

Two main records have recordedis found in transaction Operation General record and Check Point Record, which is divided into commit and update record [13].In cases of system, failure recovery is made using Check point records and Restart area contain their LSNs. So, this can be taken as stable position of the system before the start of any transaction. Transactional updates have been traced and placed in Update Record. Last transactional record has been saved in Commit record. These are later

Converted to Check Point Record . Check Point Records and General Records have the same structure and carries necessary information to performed o run do (rollback) operation [25].

Table9: Operational Record structure of $logfile

	0	1	2	3	4	5	6	7
0x00	Current LSN							
	ClientPreviousLSN(forb acklinkingofthisrecord)							
0x10	ClientUndoLSN (incaseoferrorrecovery,usuallythesameaspreviousLSN)							
	ClientDataLength (from"RedoOP",thesizeofRecord)				ClientID (owneroftheisrecord)			
0x20	RecordType (0x02forCheck PointRecord,0x01 forthegeneralRecord)				TransactionID (usedexternallybythe client (transactionmanager)togrouplogfileentries)			
	PageoverflowFlags forthisrecord (0x01-recordoverflows thepage, 0x00–nooverflow)		Padding					
0x30	Redo OP (code-checkcodesheet)		UndoOP (code-checkcodesheet)		Redo Offset (start ofRedodatafrom"Redo		RedoLength	
	UndoOffset (startofundodatafrom" RedoOP" =RedoOffset+RedoLe ngth)		UndoLength		TargetAttributeoffset (fromRedoOP,thelocati on oftarget VCN)		LCNstoFollow (0x01- Thereisthenextrecord) , 0x00-	
0x40	AttributeOffset (offsetinsideMFTrecordfo rappliedRedo/Undodata,if the changeaffectsanMFTreco rd,otherwise0x00)		OffsetwithinAttribute (theoffsetofpointapplied Redo/Undodatawithinthe attributein MFTrecordorwithin the cluster)		MFTCluster Index (Incaseofoperation for MFTrecord,thelocationof recordappliedRedo/Undo datawithinthe cluster 0000-first,0002- Second,0004-third,0006-		padding	
	TargetVCN (VCNof"$MFT"fileappliedRedo/Undodata,tobeappli edinconjunctionwith"MFTClusterIndex" MultiplyVCNwithBytes/clustertoreach correspondingbytenumberandthencountthenumber ofrecordsasmentionedinMFTClusterIndextoreach thecorrespondingMFTrecord)				padding			
0x50	TargetLCN (LCNofthediskappliedRedo/Undodata,tobeapplie dinconjunctionwith"MFTClusterIndex" MultiplyLCNwithBytes/clustertoreach correspondingbytenumberandthencountthenumber ofrecordsinthatclusterasmentionedinMFTCluster IndextoreachthecorrespondingMFTentry IncaseofbitsofBitMap,firstfourbytesareMFTRecor d Numberofthefile)				padding			
Tillthe endof data	Data							

Operations cheat sheet for Redo/Undo Codes. OperationCodeSheetforcodesmentionedin0x30-0x33in table 9 is listed in Table10 [13].

Table 10 -Redo/Undo Operation Code Sheet of $LogFile

Code	Operation	Code	Operation
0x00	Noop	0x0E	Add Index Entry Allocation
0x01	CompensationLog Record	0x0F	Delete Index Entry Allocation
0x02	Initialize File RecordSegment	0x12	Set Index EntryVCNAllocation
0x03	DeallocateFile RecordSegment	0x13	Update File Name Root
0x04	WriteEndof File RecordSegment	0x14	Update File NameAllocation
0x05	Create Attribute	0x15	Set Bits in Nonresident BitMap
0x06	Delete Attribute	0x16	Clear Bits in Nonresident BitMap
0x07	Update Resident Value	0x19	PrepareTransaction
0x08	Update NonresidentValue	0x1A	Commit Transaction
0x09	Update MappingPairs	0x1B	Forget Transaction
0x0A	Delete Dirty Clusters	0x1C	Open NonresidentAttribute
0x0B	Set NewAttributeSizes	0x1F	DirtyPage Table Dump
0x0C	Add Index Entry Root	0x20	Transaction Table Dump
0x0D	Delete Index EntryRoot	0x21	Update RecordDataRoot

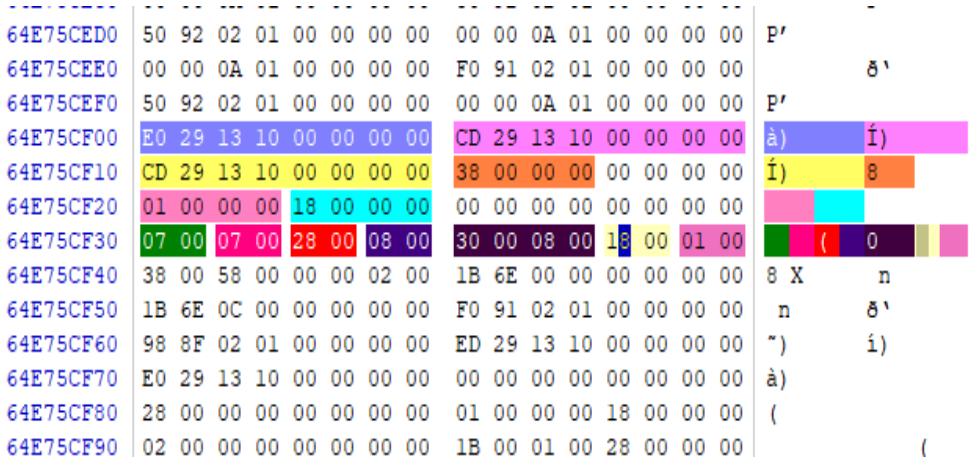


Figure 10: Example of Operation Record of $LogFile

4.2 $Recycle.Bin

$Recycle.Bin holds crucial data or evidence for the forensic investigator. Every user has his hown private bin. $Recycle.Bin folder placed inside the SID folder followed by a string which is unique for each user. Microsoft Developer Network (2009) has explained that the unique SID of each user is an alpha-numeric used to uniquely identifya user. When any user deletes a file, it will be stored in his SID folder. If there are three users are using the system will have three SID folders. One user without any permission cannot see the Bin of another user [15].

As Volume Shadow Copy generates a differential backups deleted file can be accessed by the Volume Shadow Copy if the Shadow is created before emptying the $Recycle.Bin. All activities will be maintained into the $Logfile which has also been copied into the Shadow Copy. Inorder to carryout this research it is important to understand the $Recycle. Bin structure completely.[15]

Whenever a file is deleted it produces two files:$I and $R files followed by the same six alphanumeric string. $I<abc>.<ext> contains the metadata about the deleted file and $R<abc>.<ext> holds the actual deleted file. Whenever a file is deleted its “deleted and “created” timestamps get associated with the file $Recycle.Bin is a great source of evidence for the forensic investigator.[3][15][16]

Structure of $I file is shown in Table:

Table 11:$I structure

0	8	Header(0x02)
8	16	Deleted file size
16	8	Deleted timestamp
24	4	File Name Length
28	Variable length	Filename and path

Example of $I structure of deleted folder named Account_info is shown in Fig:11

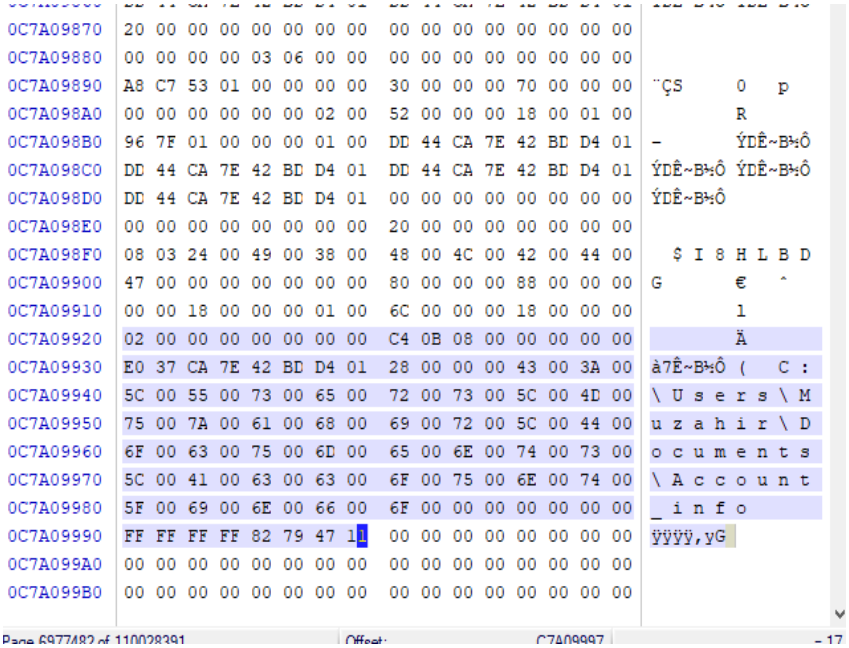


Fig 11: Example of $I file structure of Account-info folder

At location ***0x0C7A09920*** is header and its value is 2 in Windows 10. From ***0x0C7A09928-0x0C7A9935*** is ***080BC4 = 514 KB*** size of the deleted folder. From ***0C7A9930-0C7A9937*** is time stamp for deleted folder. After converting it date time stampE037CA7E42BDD401***=2/5/2019 12:05pm*** *using convertor***.** At the Offset **0x18**timestamp start of the file path described in table 12.

Table 12: Artifacts of $I file of deleted Folder Account_Info

Offset	**Remarks**	**Findings/Artifacts**
0x0C7A09920	Header	02
0x0C7A09928-35	Size	514 KB
0x0C7A09930-37	Time stamp	2/5/2019 12:05pm
0x0C7A09938-97	Path of the deleted file	C:\User\Muzahir\Documents\Account_Info

After the deletion, the two files have been generated $I and $R [15]. Both files named as $I8HLBDG for metadata of the file and $R8HBDG for the actual deleted file.

All SID folders also have index attribute files which holds the metadata of all the files contained in those folders. After analyzing the shadow copy at offset C5FE5800 I have located file ***S-1-5-21-3207469532-412700025-1003426619-1001~$I30 (90)*** that is SID of $Recycle.Bin of the user. This file contains all the deleted files and folders' metadata.[17].$I30 is an index attribute which implements the B-Tree formation which keeps a record of the deleted or overwritten files. As in MFT all the files are not deleted/removed completely in index record all the tree nodes also not deleted only marked as deleted using their $Bitmap entry. It is not guaranteed that all the files which are presented in index record are present in volume but with the help of index record file we can find in long-lasting deleted or overwritten files with their metadata which includes [17]

- ✓ Name of the file
- ✓ Parent folders
- ✓ Creation time
- ✓ Deletion time
- ✓ MFT change time
- ✓ Access time

An example of an index attribute file for SID ***S-1-5-21-3207469532-412700025-1003426619-1001~$I30***been shown in fig 11.

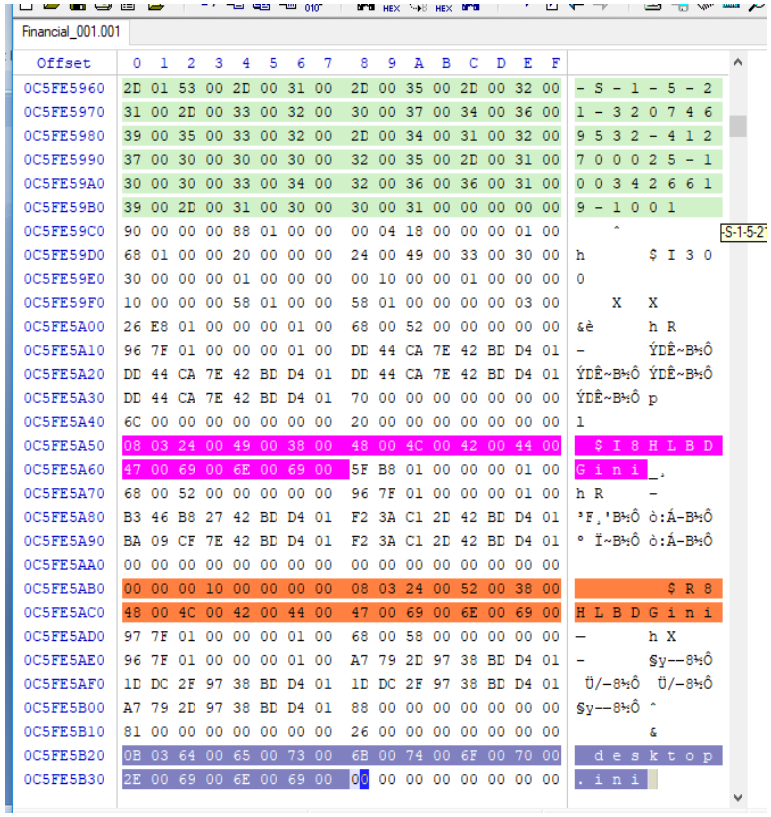


Fig12: $I30Index Attribute of the SID

Name of the file SID of the user
$I file
$R file
Desktop.ini is included in all SID folders

This chapter highlights the importance and structure of the $logfile and $Recycle.bin for the Forensic Investigator as they contain valuable information to dig out the evidence to prove in court of law.

Chapter5

Results and Analysis

In this chapter all the artifacts, findings of the experiment and results have been discussed. Volume Shadow Copy is incremental backups of 16 KB and created automatically after specific time intervals or manually whenever the user wants. By investigating Volume Shadow Copy one can view the previous version of the files. By following all the steps already discussed in Chapter II and given in the proposed forensic methodology we find the artifacts.

5.1 Composing Volume Shadow Copy from Forensic Image of Case Scenario1:

As discussed in Chapter III composing the Volume Shadow Copies is a multistep process step which includes store header, catalogs, descriptors etc. Following steps have been followed to construct a Volume Shadow Copy from forensic image for Case scenario 1(Stolen financial information) forensic image of the suspect's computer has been taken with the help of FTK imager named as **Financial_001**. Four Shadow Copies has been created for the Case1:

Volume Shadow Copy I: Contains the shadow copy of the whole system after copying the stolen information into the suspect's computer. Initially, files were saved on desktop.

Volume Shadow Copy II: The second shadow was triggered when the files were moved into the new folder account-info created on the desktop.

Volume Shadow Copy III: The third shadow was created when the folder Account_info was moved to the Document folder.

Volume Shadow Copy IV: Fourth shadow copy has been creating when the folder Account_info has been deleted

In order to traces, the changes shadow copy has been composed using the following steps.

5.1.1Presence of Volume Shadow Copy: In order to check whether the volume shadow copy is present in the system or not, catalog block offset from volume shadow header from location ***0x1E00***has been checked. After converting it into little Indian it was ***0x071F0000.*** It is the location of the first catalog block offset of the first shadow copy.

Table: Artifacts found from Volume Shadow Copy header and catalog block header

Offset	Findings/Artifacts	Remarks
0x1E00	VSS header	VSS Header always start from this location
0x71f0000	Catalog Block header	Contains the information of all shadows present in the system
0x071f0080	Entry type 0x02	Contains the metadata about the shadow copy Size ,Creation time, Shadow identifier
0x071f0100	Entry type 0x03	Contains the store offsets , store range offset ,store bitmap offset

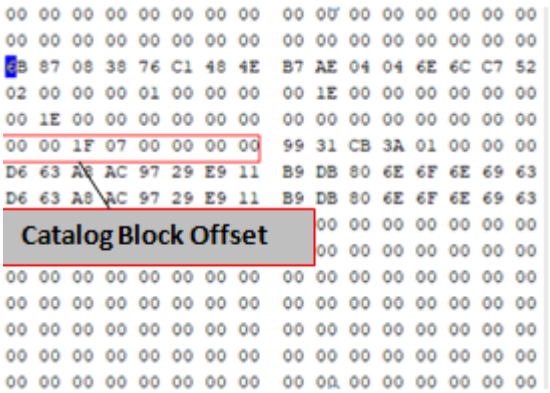


Fig13: Volume Shadow Copy Headers

By accessing the Catalog Block from the above-mentioned Catalog Block header has been decoded as shown in fig 14. There are total 8 catalog entries against 4 shadow copies. Entry type 0x02 contains the metadata of the Volume Shadow Copy and entry type 0x03 contains data about the actual volume snapshots shown in fig 14. These two entry types are decoded to according to the Table 4 and Table5 described in chapter II.

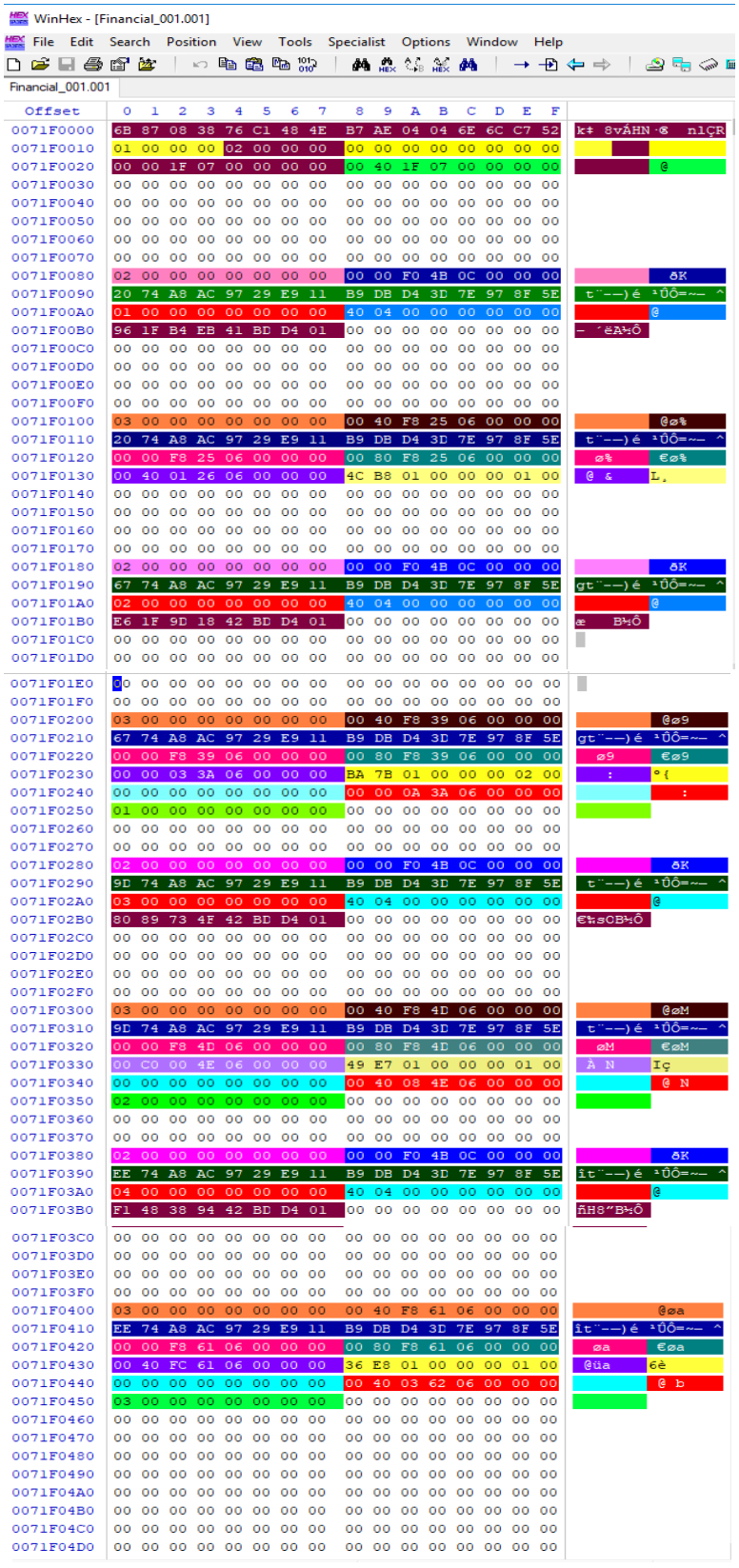


Fig14: Catalog Block entries type 0x02 and 0x03

The offset of the first store block has been found and decoded from catalog block header as shown in fig 15.

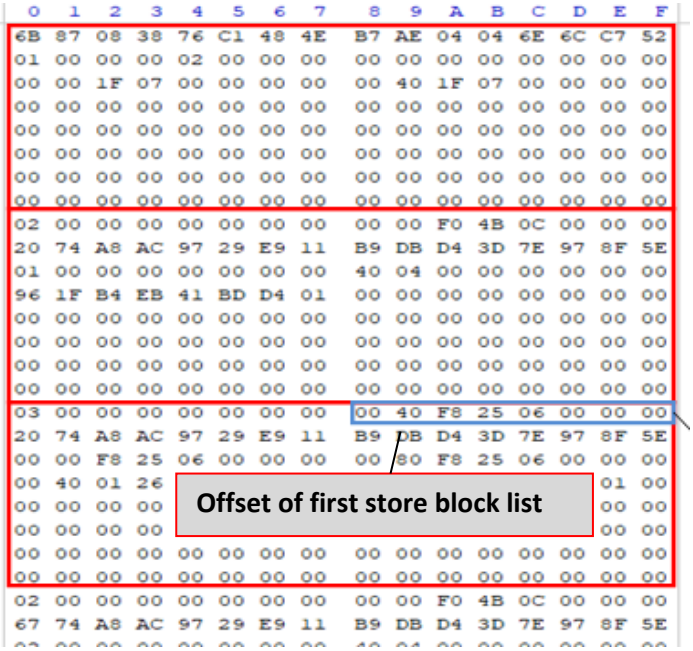


Fig 15: Catalog block header and catalog entry types

5.1.2Locating Catalog Store Block List

From entry type 0x03 we locate the offset of first store block list. After converting it into little Endean we reached first store block list located at **0x625F84000** shown in fig16.

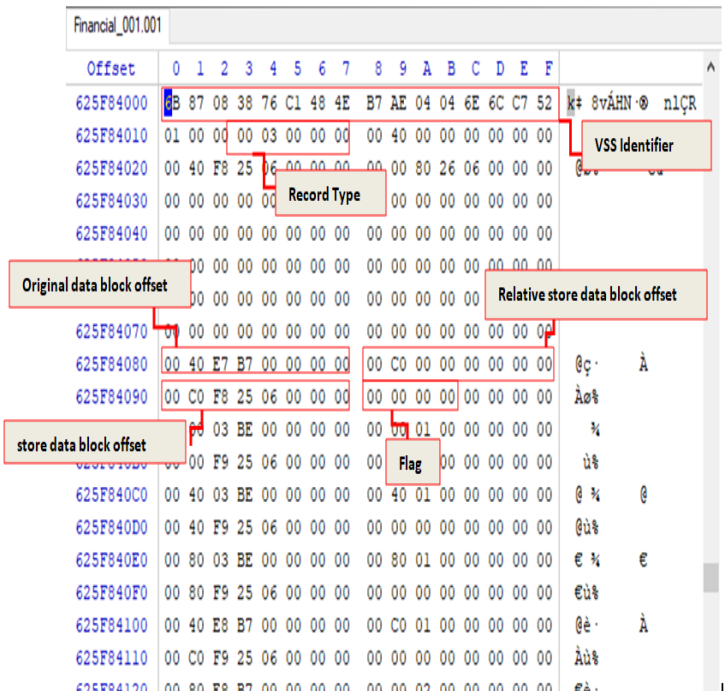


Fig 16: Shadow Store block list of first Shadow

Table :Artifacts found from store block list

Offset	**Findings/artifacts**	**Remarks**
0x625f84000	VSS Identifier	Volume shadow copy name
0x625f84014	Record type	Store block list
0x625f84080	Original data block offset	Data block of active volume
0x625f84090	Store data block offset	Offset of first volume shadow copy
0x625f84098	Flag value	Descriptor value 0=Normal descriptor

From store block list header type 0x03 offset of first store block is found ***0x0625f8000*** as shown in fig17.

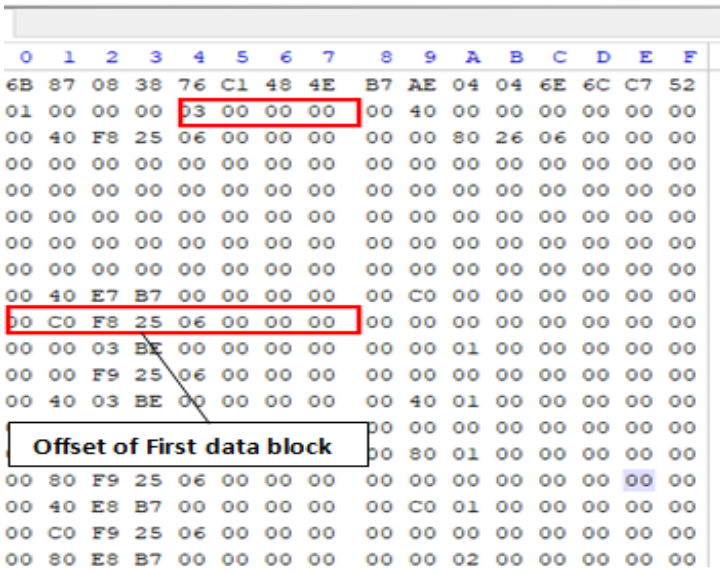


Fig 17: Shadow Store block list header

Shadow block list header is of type 0x03 and contains the store data block offset and flag descriptor value which is 0 indicates normal block descriptor.

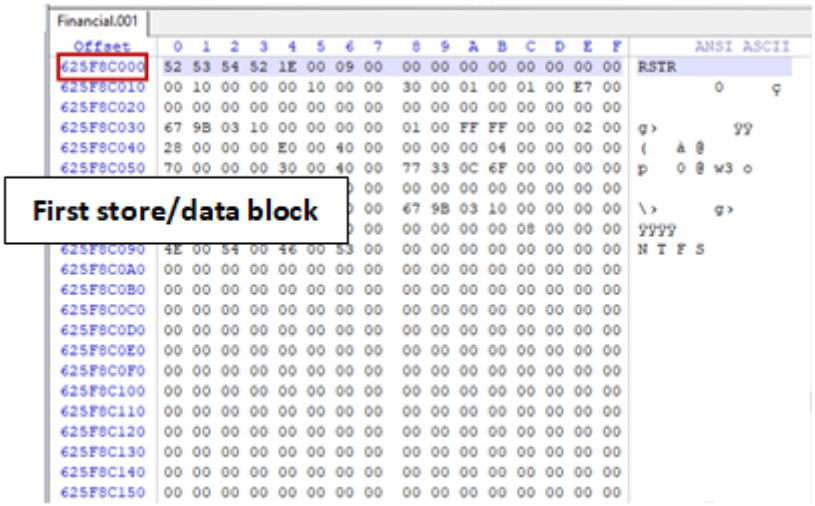
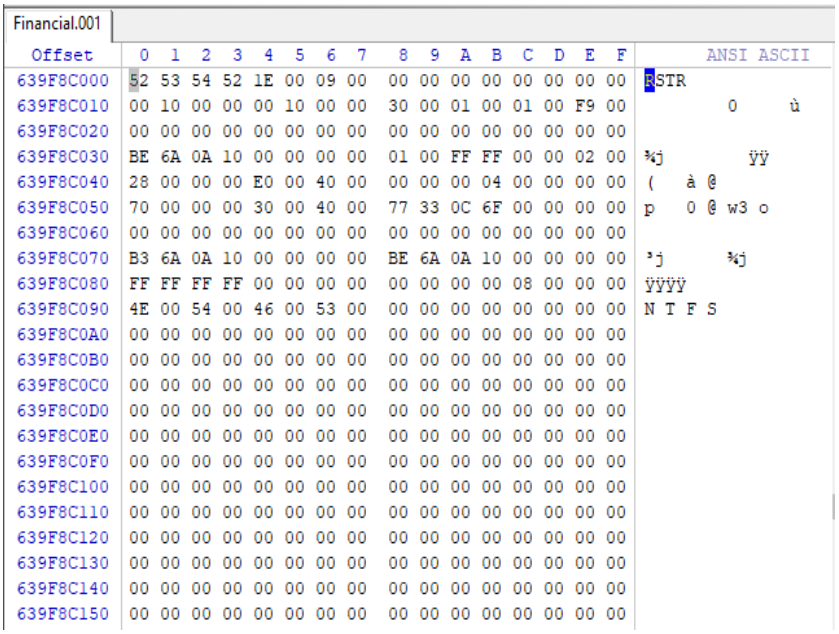


Fig 18: Snapshot of First data block

First data store indicating $logfile captured in 1st Volume Shadow Copy. It only saves the changes to the original $logfile. Hence the traces can also be traced from $logfile as well.

Second data store found at the location of ***0x639f84000***as shown in fig19.

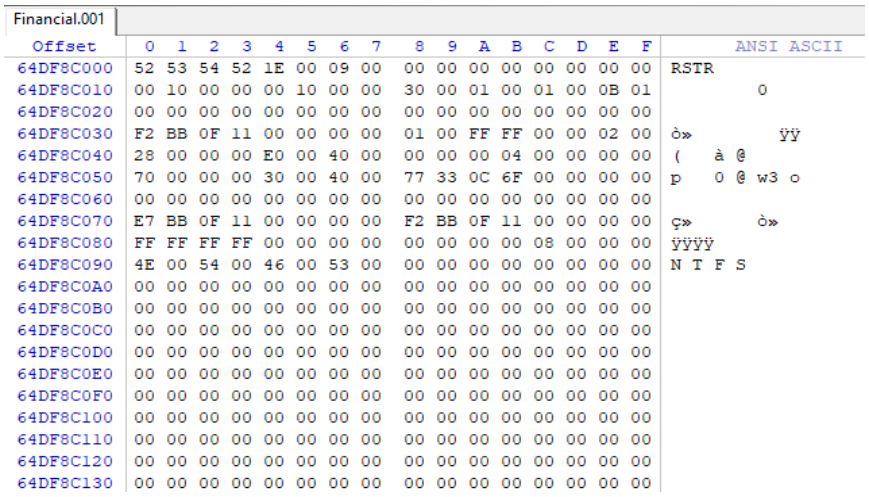


Financial.001

Offset	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F	ANSI ASCII
639F8C000	52	53	54	52	1E	00	09	00	00	00	00	00	00	00	00	00	RSTR
639F8C010	00	10	00	00	00	10	00	00	30	00	01	00	01	00	F9	00	0 ù
639F8C020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C030	BE	6A	0A	10	00	00	00	00	01	00	FF	FF	00	00	02	00	¾j ÿÿ
639F8C040	28	00	00	00	E0	00	40	00	00	00	00	04	00	00	00	00	(à @
639F8C050	70	00	00	00	30	00	40	00	77	33	0C	6F	00	00	00	00	p 0 @ w3 o
639F8C060	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C070	B3	6A	0A	10	00	00	00	00	BE	6A	0A	10	00	00	00	00	³j ¾j
639F8C080	FF	FF	FF	FF	00	00	00	00	00	00	00	00	08	00	00	00	ÿÿÿÿ
639F8C090	4E	00	54	00	46	00	53	00	00	00	00	00	00	00	00	00	N T F S
639F8C0A0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C0B0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C0C0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C0D0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C0E0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C0F0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C100	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C110	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C120	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C130	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C140	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
639F8C150	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	

Fig: 19 Snapshot of Second Store: starts with $logfile

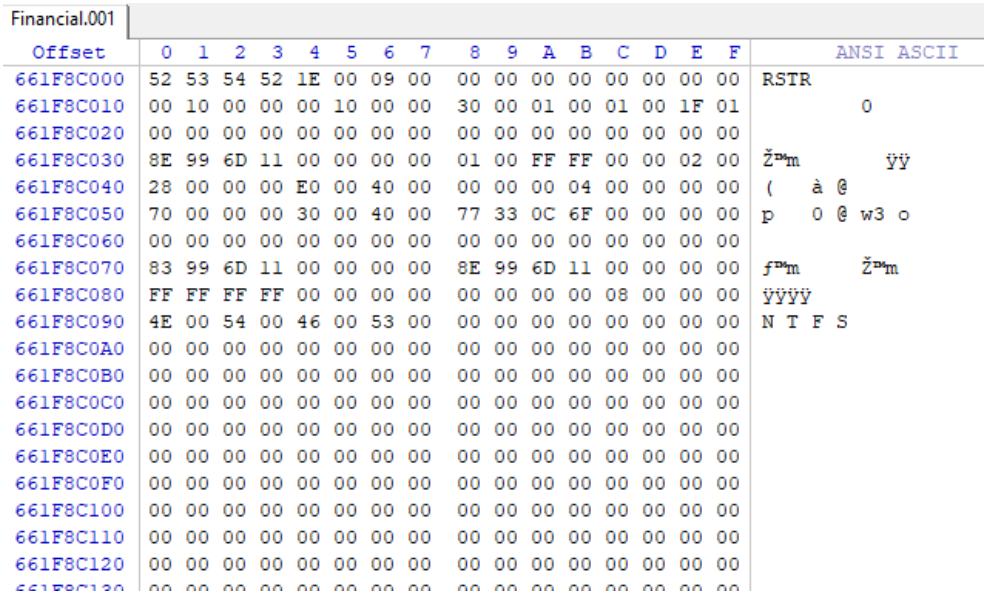
In the same way, third and fourth data stores have been traced from catalog block header entry type 0x03 and it is observed that theses all blocks start with $logfile shown in fig 20 and 21.



Financial.001

Offset	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F	ANSI ASCII
64DF8C000	52	53	54	52	1E	00	09	00	00	00	00	00	00	00	00	00	RSTR
64DF8C010	00	10	00	00	00	10	00	00	30	00	01	00	01	00	0B	01	0
64DF8C020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C030	F2	BB	0F	11	00	00	00	00	01	00	FF	FF	00	00	02	00	ò» ÿÿ
64DF8C040	28	00	00	00	E0	00	40	00	00	00	00	04	00	00	00	00	(à @
64DF8C050	70	00	00	00	30	00	40	00	77	33	0C	6F	00	00	00	00	p 0 @ w3 o
64DF8C060	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C070	E7	BB	0F	11	00	00	00	00	F2	BB	0F	11	00	00	00	00	ç» ò»
64DF8C080	FF	FF	FF	FF	00	00	00	00	00	00	00	00	08	00	00	00	ÿÿÿÿ
64DF8C090	4E	00	54	00	46	00	53	00	00	00	00	00	00	00	00	00	N T F S
64DF8C0A0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C0B0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C0C0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C0D0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C0E0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C0F0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C100	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C110	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C120	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
64DF8C130	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	

Fig:20 Snapshot of Third Store



Financial.001

Offset	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F	ANSI ASCII
661F8C000	52	53	54	52	1E	00	09	00	00	00	00	00	00	00	00	00	RSTR
661F8C010	00	10	00	00	00	10	00	00	30	00	01	00	01	00	1F	01	0
661F8C020	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C030	8E	99	6D	11	00	00	00	00	01	00	FF	FF	00	00	02	00	Ž™m ÿÿ
661F8C040	28	00	00	00	E0	00	40	00	00	00	00	04	00	00	00	00	(à @
661F8C050	70	00	00	00	30	00	40	00	77	33	0C	6F	00	00	00	00	p 0 @ w3 o
661F8C060	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C070	83	99	6D	11	00	00	00	00	8E	99	6D	11	00	00	00	00	ƒ™m Ž™m
661F8C080	FF	FF	FF	FF	00	00	00	00	00	00	00	00	08	00	00	00	ÿÿÿÿ
661F8C090	4E	00	54	00	46	00	53	00	00	00	00	00	00	00	00	00	N T F S
661F8C0A0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C0B0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C0C0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C0D0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C0E0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C0F0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C100	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C110	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	
661F8C120	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	

Fig:21 Snapshot of Fourth Store

The Data stores are manipulated in a way that the changes traced from the latest store first and then to the next latest and the oldest snapshot will be manipulated at the end. Every 16 KB block has to be manipulated to verify the modification of the original contents. Flag in index record was set as 0x00 in all four shadow copies which indicate normal block.

In case of normal block whole block is mapped to the original block to get the data. In order to generate a Volume Shadow Copy of the forensic investigator has to check all 16 KB blocks. If 16 KB block is not modified, then read data from the original position. If the index record of the block does not exist, then read the corresponding bit from the bitmap.[8]

5.2 Reading snapshot from store data:

BD=Block Descriptor *CV=Current Volume*

FD=Forward block descriptor *RBD=Reverse Block Descriptor*

```
If BD=1 then
    If BD=Over_Lay then
    Use overlay block descriptors
    If FD=1 and Next_store=1 then
    Read data from store   *Using relative store offset
    Else
    Read data from the original block
  Else
    IF next_store =1 then
    Read data from next store
    Else
    IF RBD =1 then
    Read data from CV
    Else If Active_store=most recent
    Read data from original Volume
```

5.3 Analysis of $Recycle.Bin

Whenever a file is deleted its two copies have been generated into the $Recycle.Bin. $R file and $I file with the same random integer values. If a file named hello.txt is deleted its two files $Igh4563 and $RIgh4563 files will be generated. $R contains the actual data of the file that has been deleted and $I file contains the metadata of the deleted file. The very first problem that a forensic investigator has to face is a large amount of data to be processed as 2 files are generated in case of deletion of 1 file [3].

After analyzing the shadow copy at offset ***0xC5FE5800*** I have located file ***S-1-5-21-3207469532-412700025-1003426619-1001~$I30 (90)*** that is SID of $Recycle.Bin ofthe user. This file contains the all the deleted files and folders' metadata.$I30 is an index attribute which implements the B-Tree formation which keeps a record of the deleted or overwritten files.AS in MFT all the files are not deleted/removed completely in index record all the tree nodes also not deleted only marked as deleted using their $Bitmap entry. Attribute Index has shown in fig 22.

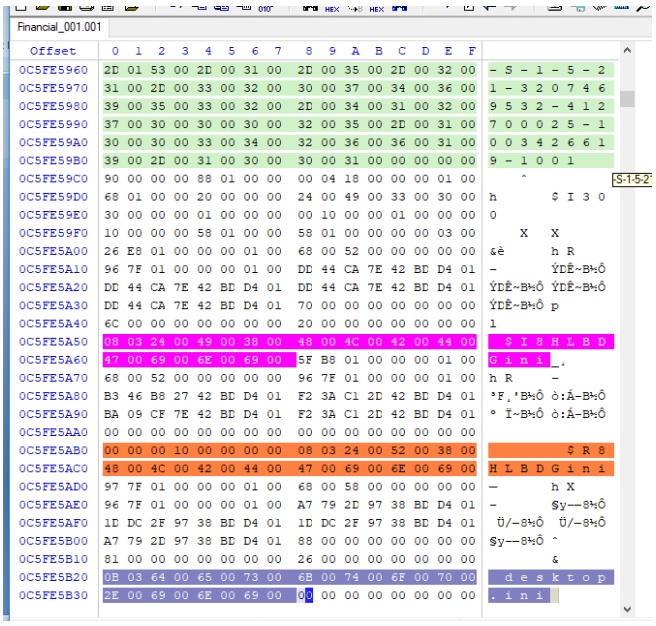


Fig:22 $I 30 index attributes of the file

Name of the file SID of the user
$I file
$R file
Desktop.ini is included in all SID folders

Structure of the $I file is shown in table.12

Table.12 Structure of $I file

0	8	Header(0x02)
8	16	Deleted file size
16	8	Deleted timestamp
24	4	File Name Length
28	Variable length	Filename and path

n NTFS file system each file has an MFT table entry. The file which was deleted from the system actually it is not deleted only its entry in MFT table marked as deleted actual contents remain there until or unless its overwritten [21]. But if the contents of the files are overwritten then it cannot be recovered. Whenever any file deleted accidentally or by choice its contents are saved before deletion into the Volume Shadow Copy. So, the file can be recovered using Volume Shadow Copy. First Volume Shadow Copy is the complete copy. From Second to onward only changes have been copied. In our experiment, we have changed the location of the stolen files and then at the end, they were deleted so in shadows only the differences have been saved.

The second shadow shows that the Account_info folder is placed at the desktop shown in fig 23.

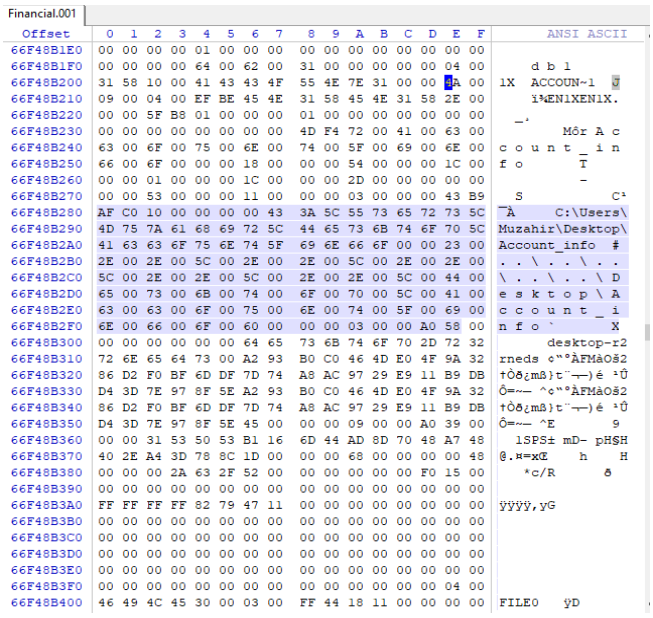


Fig:23Account_info in 2nd in Shadow copy

After the generation of second shadow copy folder was moved on the document folder and shown in fig 24 and 25.

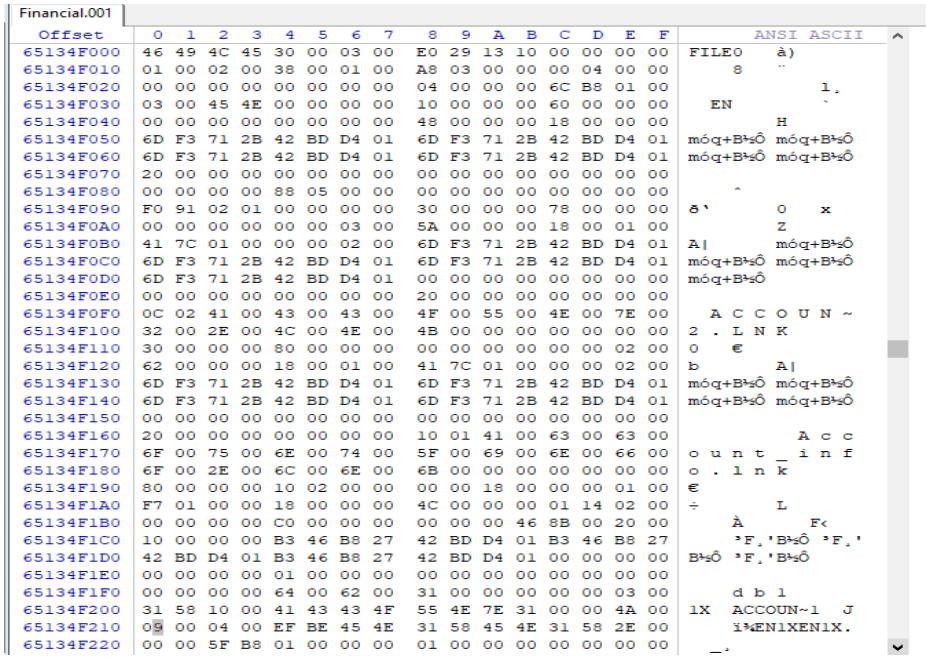


Fig 24: Account_info file in a 4th shadow copy

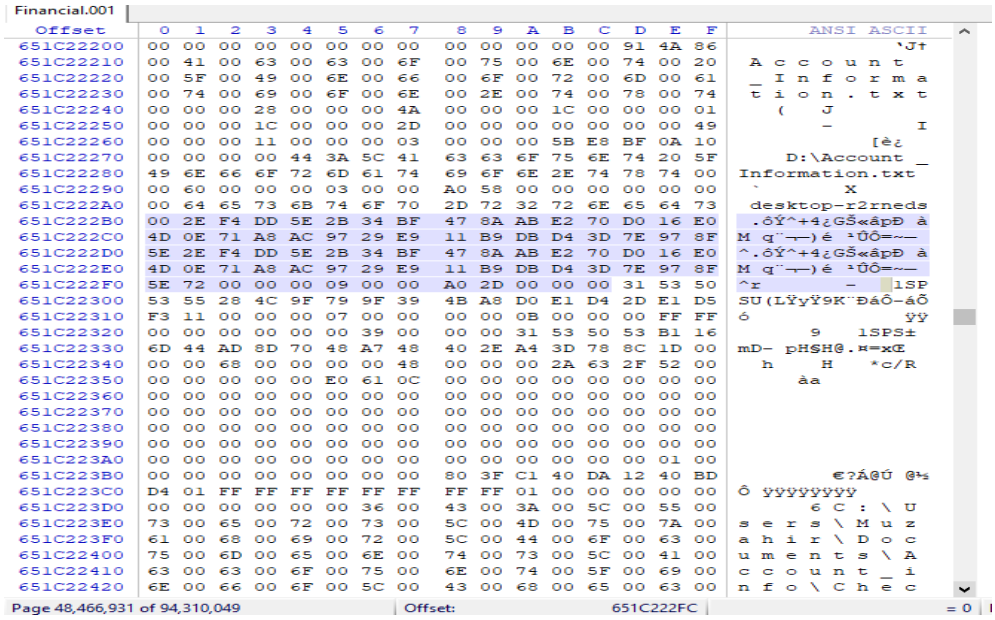


Fig 25: Account_information.txt file in 3rd shadow copy is found D director

In fourth Shadow Copy file, Check-bank.jpg file has been moved from desktop to the Document folder, so in 4 the Volume Shadow Copy only the change which is the path of the file saved. The change path of the file is shown in 26.

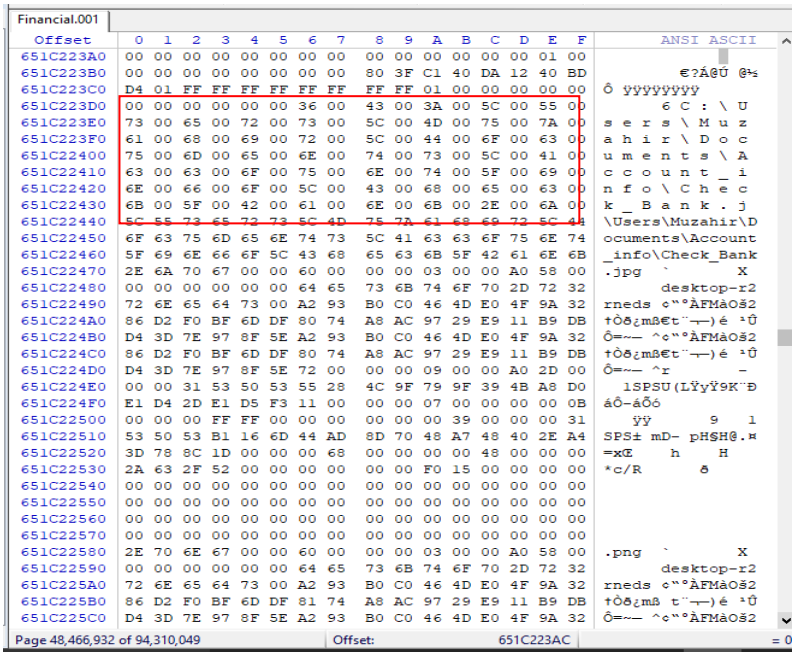


Fig: 26 Check_bank in 4 the shadow copy

Bank_statement file was deleted and its deletion through recycle.bin is saved in 4th shadow copy as shown in fig 27:

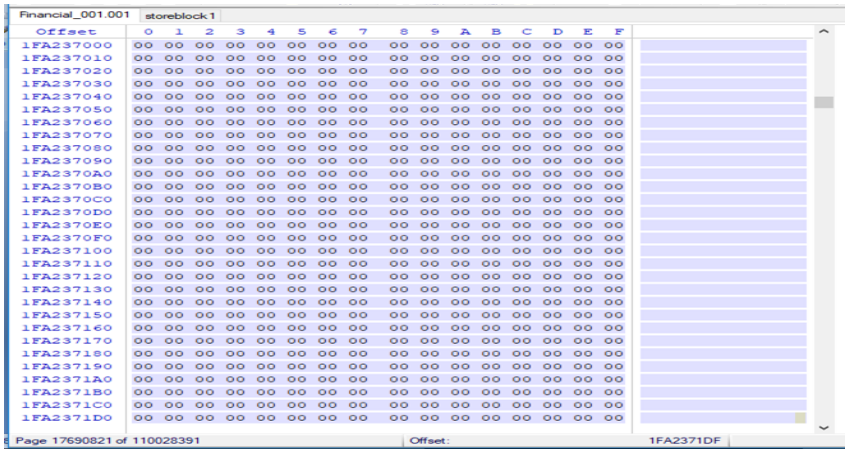


Fig.27: Deleted file Bank_ statement.jpg

The file Bak_statement is recovered back by mapping the shadow copy to the original data block and successfully recover shown in fig 29.

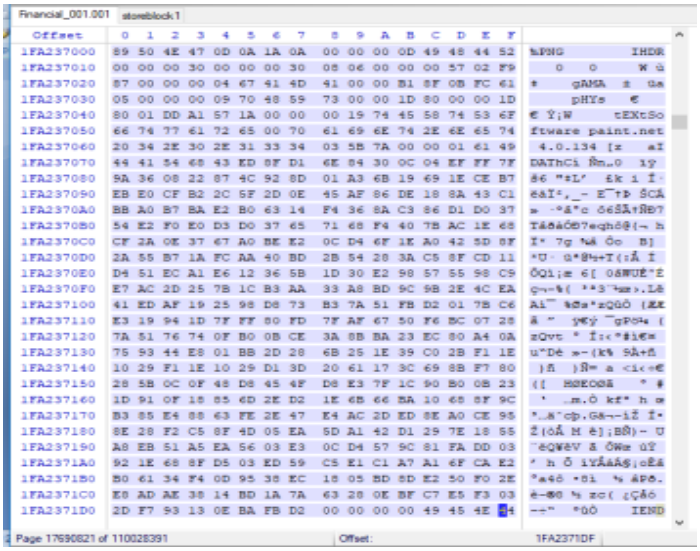


Fig.29: Bank _statement.jpg file is recovered

It might possible that all the artifacts related to the deleted file is no longer present

into the MFT but after doing the analysis of $log file examiner can dig out the important artifacts.

Every shadow copy or store starts with the $logfile. As the first store is placed at offset ***0x625f84000 .*** From index record store's relative offset and original file offset pointing to the $log file from original volume and shadow copy found which is shown in fig 29.

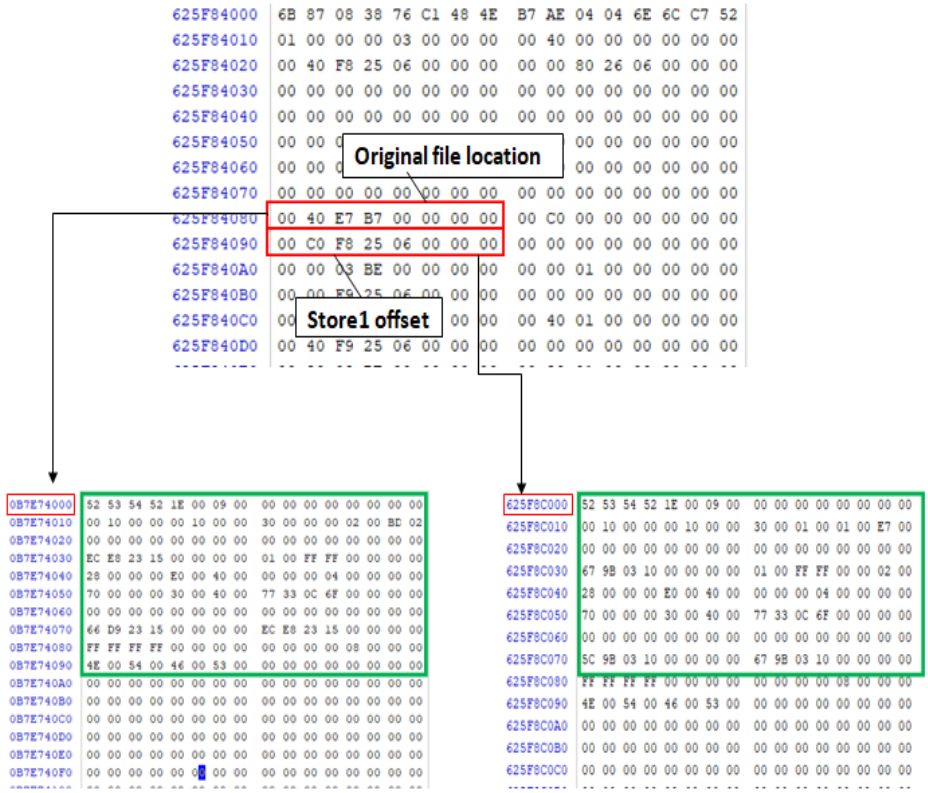


Fig 29: Original volume offset, and First Volume Shadow copy Offset

Table:16 Artifacts from the first catalog block list

Offset		Remarks	Findings/Artifacts
625F84080-87	8 Bytes	Original file location	E7B74000 (Original $Logfile Location)
625F84090-97	8Bytes	Location of the Store1(1st Shadow copy)	625F8C000 (Offset of the 1st shadow Copy)

The original file location is a location of the $log file of the Active volume and store1's offset represents the start of the Volume snapshot which contains $log file values which save the state of the $logfile for that specific point in time when shadow copy was created. It's important from the forensic point of view that investigator can investigate from that $logfile which saves the state of the system.

Every file has LSN and can be identified with its LSN in a $log file. Folder Account_info has LSN ***0E72AC23*** in its $MFT entry. In $Log file LSN of folder Account_info which was deleted and contains the all three files 1) Bank_Staement.jpg, Bankcheck.jpg,and Account_information.txt have been found at the location ***0x64E75CDE0***.Corelation between $MFT and $log file for the said folder has been shown in fig 30.

Fig 30: LSN of a file in $MFT and $Logfile of Volume Shadow Copy

5.4 File Deletion case.

Two types of file deletion have been supported by NTFS: using Recycle Bin and bypassing the Recycle bin.

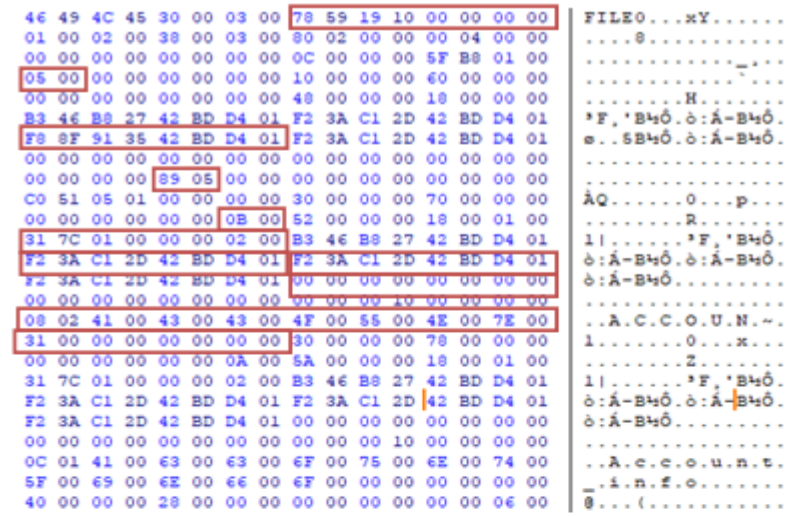
Deleting files using Recycle Bin. As soon as the file is deleted it is moved into the recycle bin after renaming. Changes also made in $MFT file records and $logfile transactional records.

Changes in MFT Record of File is moved to Recycle.Bin after Deletion

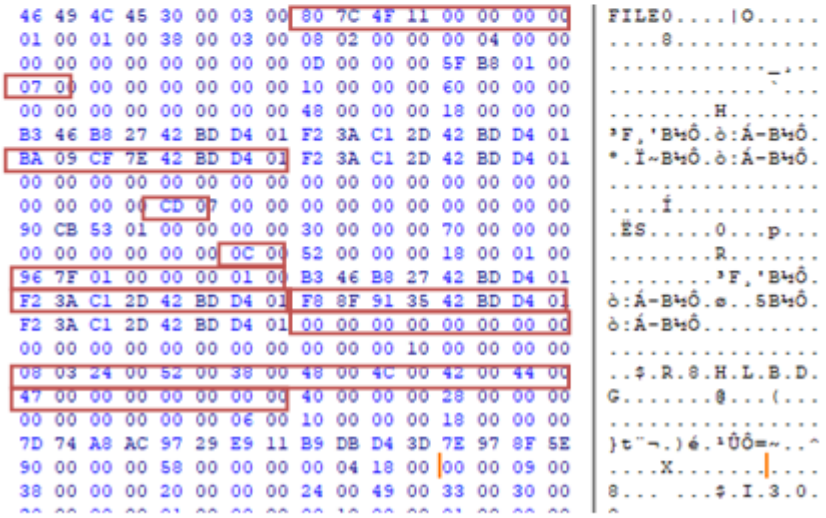
- Log file Sequence Number is updated at offset0x08-0x0F.
- Next attribute ID at offset0x28 is incremented by 1, as previous "next attribute ID" has been assigned to attribute 0x30 as it has been changed.
- Updatesequencenumberatoffset0x30-0x31,0x1FE-0x1FFand0x3FE-0x3FFare updated.
- MFT update time at offset 0x60 is updated in attribute 0x10 $Standard_Information.
- Security ID is updated at offset 0x84 is updated in attribute 0x10 $Standard_Information.

- As the name of the file is changed, so attribute length is changed at offset 0x9C-0x9 -0x30 $File_Name.
- Attribute ID at 0xA6-0xA7of attribute 0x30 $File_Name is incremented to "Next attribute ID" mentioned in the original record, as this attribute has been updated and will be accordingly mentioned in $logFile to differentiate the updated attribute 0x30 from the previous record.
- The file is indexed under the current user folder in recycle bin instead of the folder under which it was previously held and accordingly "File reference to the parent directory" at offset 0xB0-0xB7 is updated.
- File modified time, and MFT update time are updated at offset 0xC0-0xC7 and 0xC8-0xC Fin 0x30$File_Name, respectively.
- Allocated and real size (of payload of attribute 0x80 $Data) is updated at offset 0xD8-0xDFand 0xE0-0xE7 in 0x30 $File_Name, respectively.
- File name length, file name namespace and file name are updated starting from attribute0xF0in0x30 $File_Name.

The comparative screenshot is shown in Figure 32. Most of the above- mentioned offsets are specific to following screenshot for ease of understanding.



Folder before deletion



Folder after deletion

Figure32: Comparative Screen Shot of MFT for File Deletion

Changes in MFT Record of New Parent Folder (Current User Folder in RecycleBin)

When the **file is moved to RecycleBin** . MFT record of parent folder under which the file after deletion has been indexed (moved) (shown at offset 0xB0-0xB7 in Figure32) is shown in Figure 33.

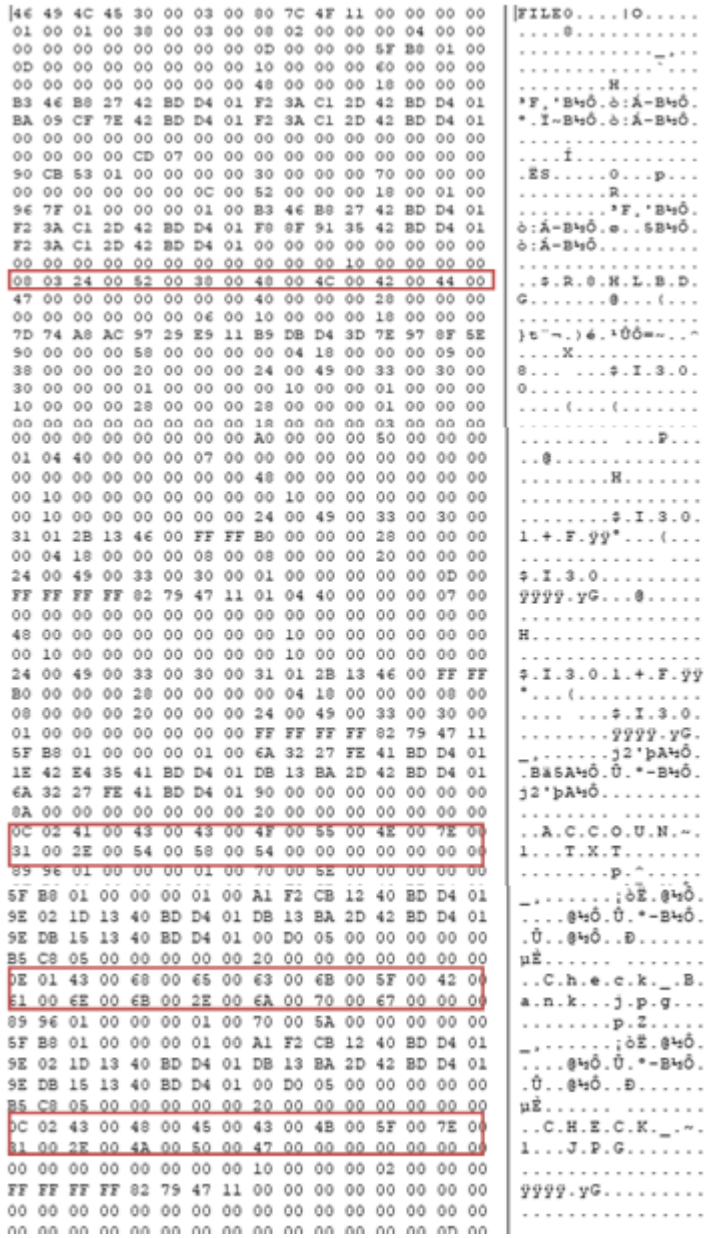


Figure 33-ExampleofCurrentUserFolder of Recycle MFTRecord

This MFT record is, in fact, a folder of current user in recyclebin folder referred at

Offset 0xB0-0xB7, shown inFigure34.

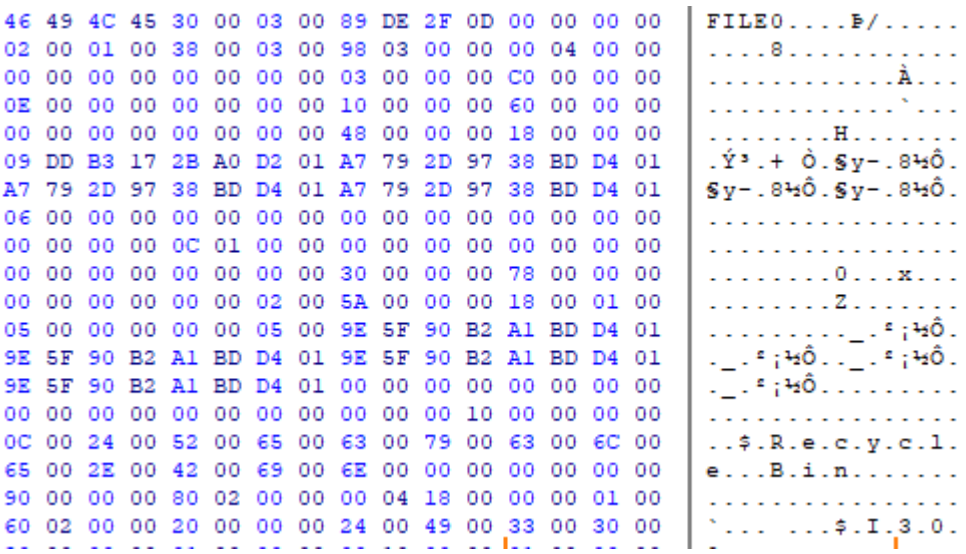


Figure 34: Example of Recycle Bin MFT Record

When any file or folder has been deleted two pair of files are generated as discussed in chapter 2.$I and $R with random characters the same for both files.$I contains the metadata and $R is an actual deleted file[22].The interesting thing is when the folder is deleted its pair of files have been generated with new names but the contents of the folder remain the same with their original names.

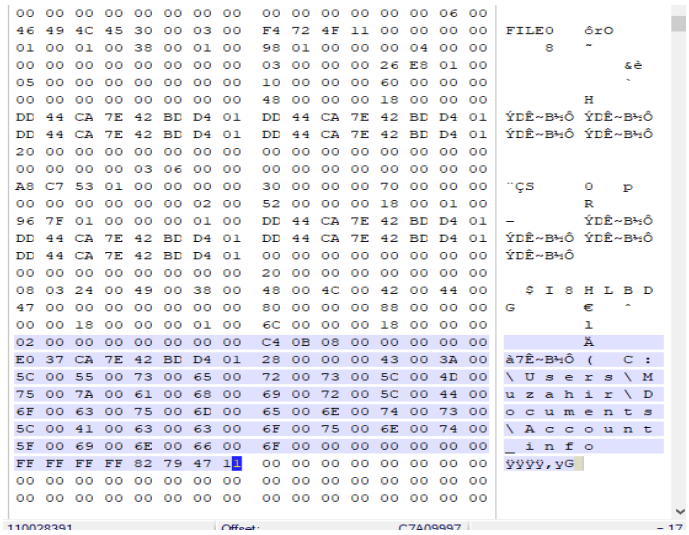


Figure 35 -$I8HLBDG MFTRecord Structure

Transactions Logged in $LogFile when File is pushed to RecycleBin. Deletion process initiates an operational records' long chain which is recorded in $LogFile. This chain starts from logging the transaction of allocating MFT record in $MFT for newly created $Ixxx file and terminates at updating index entry of root record 05 of $I. This chain contains numerous small subchain each representing a sequence of transactions. This chain contains numerous small sub chains each representing a sequence of transactions. Important transactions logged in $logfile(in sequential order) are discussed below.

SubChain1–MFTRecordGenerationof$Ixxx

- A new file record($Ixxx) generated for Current User Folder of Recycle Bin to point to MFT location of deleted File
- New File($Ixxx) indexed in Current User Folder of RecycleBin
- Initialization of MFT Record of $Ixxx

SubChain2–UpdateRecycleBin

- Update MFT of RecycleBin
- UpdateIndexEntryof Current User in RecycleBin

SubChain3. Update 0x80 of MFT record of $Ixxx.

Sub Chain 4. Update Index Entry of $Ixxx in Current User Folder of Recycle Bin.

SubChain5–Renaming Deleted File and Updating Associated Records

- Removing Deleted file entry(Long Filename) from Parent Folder
- Removing Original File Name from MFT record of Deleted File
- Removing Deleted file entry(Short File name) from Parent Folder
- Renaming Deleted File to $Rxxx in MFT Record of Deleted File
- Clearing Index Entries from Recycle Bin of Current User Folder

A new Cluster Allocated for Non-Resident INDX of Current User Folder of RecycleBin in $BitMap

- Creating Attr. 0xA0 in MFT in RecycleBin of Current User Folder
- Creating Attr. 0xB0 in MFT in RecycleBin of Current User Folder
- Updating Attr. 0xB0 inMFT in RecycleBin of Current User Folder
- Updating Attr. 0xA0 in MFT in RecycleBin of Current User Folder
- Add index entries less Renamed Deleted File $Rxxxin MFT in RecycleBin of Current User Folder.
- Creating Attr. 0x90 in MFT in RecycleBin of Current User Folder
- Add index entry of renamed file $Rxxxto INDX of Attr.0xA0

SubChain6–UpdatingRecycleBin

- Update MFT Record in RecycleBin of Current User Folder
- Updating MFT Record in RecycleBin of Current User Folder

SubChain7.

- Updat MFT Record of Renamed Deleted File $Rxxx (Attr.0x10 only Security ID, Quota charged an Update Sequence Number).

Sub Chain 9 – Updating MFT Records in Chain of Parent Folders

- Update MFT Record of Parent Folder
- Update MFT Record of Parent of the Parent Folder

Example of operational record has been shown in fig 36.

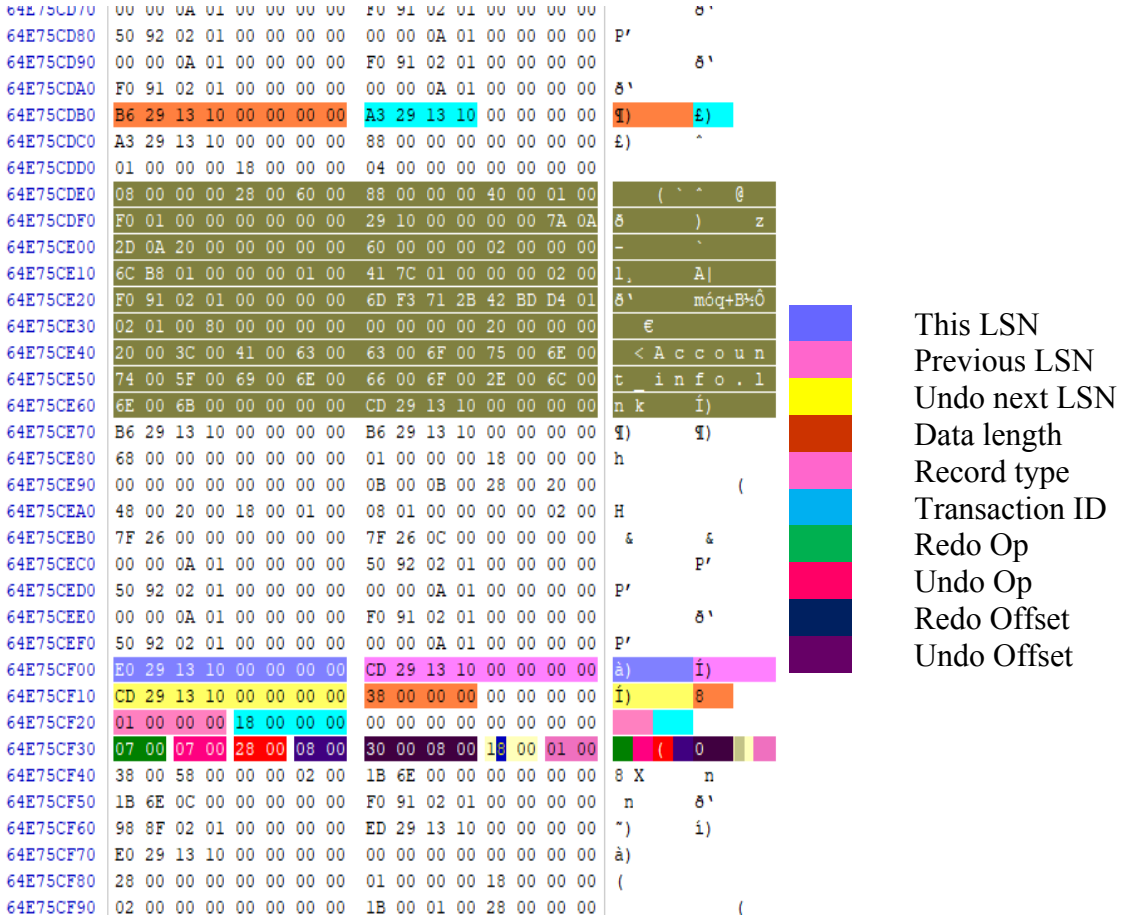
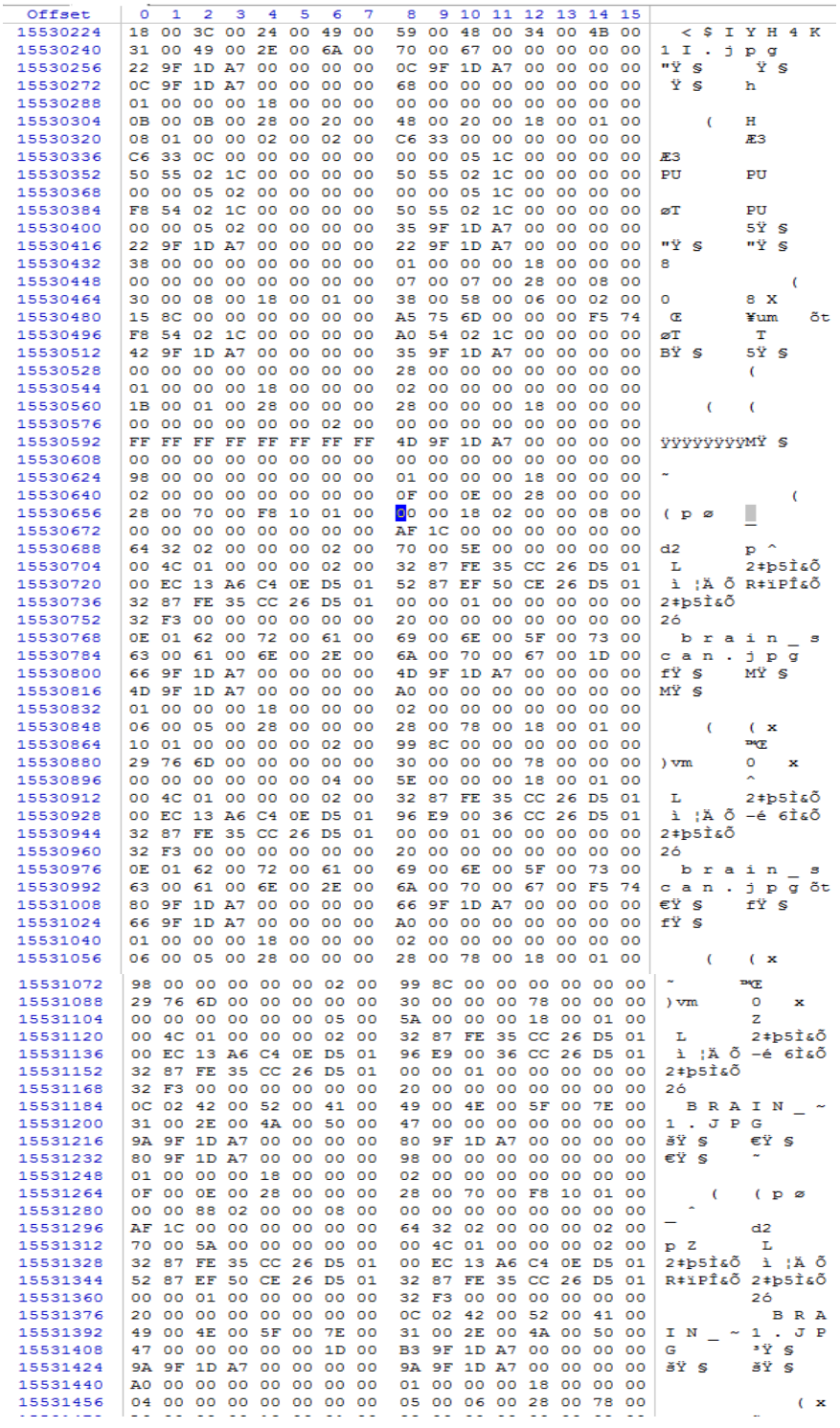


Fig 36: Operational Record of folder Account_info

At location, ***0x64E5CF20*** contains the Record type which is ***0x01***showing that it's a general record. Transaction ID has 0x18 page overflow flag and its value at location 0x 64E5CF28 is 0x00 which indicate that page is not overflowing. In the same way, remaining operational record has been.

$logfile data facts of Blood_count.docx represent the traces of the all the transactions done with the file as shown in fig 37



Offset	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
15530224	18	00	3C	00	24	00	49	00	59	00	48	00	34	00	4B	00
15530240	31	00	49	00	2E	00	6A	00	70	00	67	00	00	00	00	00
15530256	22	9F	1D	A7	00	00	00	00	0C	9F	1D	A7	00	00	00	00
15530272	0C	9F	1D	A7	00	00	00	00	68	00	00	00	00	00	00	00
15530288	01	00	00	00	18	00	00	00	00	00	00	00	00	00	00	00
15530304	0B	00	0B	00	28	00	20	00	48	00	20	00	18	00	01	00
15530320	08	01	00	00	02	00	02	00	C6	33	00	00	00	00	00	00
15530336	C6	33	0C	00	00	00	00	00	00	00	05	1C	00	00	00	00
15530352	50	55	02	1C	00	00	00	00	50	55	02	1C	00	00	00	00
15530368	00	00	05	02	00	00	00	00	00	00	05	1C	00	00	00	00
15530384	F8	54	02	1C	00	00	00	00	50	55	02	1C	00	00	00	00
15530400	00	00	05	02	00	00	00	00	35	9F	1D	A7	00	00	00	00
15530416	22	9F	1D	A7	00	00	00	00	22	9F	1D	A7	00	00	00	00
15530432	38	00	00	00	00	00	00	00	01	00	00	00	18	00	00	00
15530448	00	00	00	00	00	00	00	00	07	00	07	00	28	00	08	00
15530464	30	00	08	00	18	00	01	00	38	00	58	00	06	00	02	00
15530480	15	8C	00	00	00	00	00	00	A5	75	6D	00	00	00	F5	74
15530496	F8	54	02	1C	00	00	00	00	A0	54	02	1C	00	00	00	00
15530512	42	9F	1D	A7	00	00	00	00	35	9F	1D	A7	00	00	00	00
15530528	00	00	00	00	00	00	00	00	28	00	00	00	00	00	00	00
15530544	01	00	00	00	18	00	00	00	02	00	00	00	00	00	00	00
15530560	1B	00	01	00	28	00	00	00	28	00	00	00	18	00	00	00
15530576	00	00	00	00	00	00	02	00	00	00	00	00	00	00	00	00
15530592	FF	FF	FF	FF	FF	FF	FF	FF	4D	9F	1D	A7	00	00	00	00
15530608	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
15530624	98	00	00	00	00	00	00	00	01	00	00	00	18	00	00	00
15530640	02	00	00	00	00	00	00	00	0F	00	0E	00	28	00	00	00
15530656	28	00	70	00	F8	10	01	00	00	00	18	02	00	00	08	00
15530672	00	00	00	00	00	00	00	00	AF	1C	00	00	00	00	00	00
15530688	64	32	02	00	00	00	02	00	70	00	5E	00	00	00	00	00
15530704	00	4C	01	00	00	00	02	00	32	87	FE	35	CC	26	D5	01
15530720	00	EC	13	A6	C4	0E	D5	01	52	87	EF	50	CE	26	D5	01
15530736	32	87	FE	35	CC	26	D5	01	00	00	01	00	00	00	00	00
15530752	32	F3	00	00	00	00	00	00	20	00	00	00	00	00	00	00
15530768	0E	01	62	00	72	00	61	00	69	00	6E	00	5F	00	73	00
15530784	63	00	61	00	6E	00	2E	00	6A	00	70	00	67	00	1D	00
15530800	66	9F	1D	A7	00	00	00	00	4D	9F	1D	A7	00	00	00	00
15530816	4D	9F	1D	A7	00	00	00	00	A0	00	00	00	00	00	00	00
15530832	01	00	00	00	18	00	00	00	02	00	00	00	00	00	00	00
15530848	06	00	05	00	28	00	00	00	28	00	78	00	18	00	01	00
15530864	10	01	00	00	00	00	02	00	99	8C	00	00	00	00	00	00
15530880	29	76	6D	00	00	00	00	00	30	00	00	00	78	00	00	00
15530896	00	00	00	00	00	00	04	00	5E	00	00	00	18	00	01	00
15530912	00	4C	01	00	00	00	02	00	32	87	FE	35	CC	26	D5	01
15530928	00	EC	13	A6	C4	0E	D5	01	96	E9	00	36	CC	26	D5	01
15530944	32	87	FE	35	CC	26	D5	01	00	00	01	00	00	00	00	00
15530960	32	F3	00	00	00	00	00	00	20	00	00	00	00	00	00	00
15530976	0E	01	62	00	72	00	61	00	69	00	6E	00	5F	00	73	00
15530992	63	00	61	00	6E	00	2E	00	6A	00	70	00	67	00	F5	74
15531008	80	9F	1D	A7	00	00	00	00	66	9F	1D	A7	00	00	00	00
15531024	66	9F	1D	A7	00	00	00	00	A0	00	00	00	00	00	00	00
15531040	01	00	00	00	18	00	00	00	02	00	00	00	00	00	00	00
15531056	06	00	05	00	28	00	00	00	28	00	78	00	18	00	01	00
15531072	98	00	00	00	00	00	02	00	99	8C	00	00	00	00	00	00
15531088	29	76	6D	00	00	00	00	00	30	00	00	00	78	00	00	00
15531104	00	00	00	00	00	00	05	00	5A	00	00	00	18	00	01	00
15531120	00	4C	01	00	00	00	02	00	32	87	FE	35	CC	26	D5	01
15531136	00	EC	13	A6	C4	0E	D5	01	96	E9	00	36	CC	26	D5	01
15531152	32	87	FE	35	CC	26	D5	01	00	00	01	00	00	00	00	00
15531168	32	F3	00	00	00	00	00	00	20	00	00	00	00	00	00	00
15531184	0C	02	42	00	52	00	41	00	49	00	4E	00	5F	00	7E	00
15531200	31	00	2E	00	4A	00	50	00	47	00	00	00	00	00	00	00
15531216	9A	9F	1D	A7	00	00	00	00	80	9F	1D	A7	00	00	00	00
15531232	80	9F	1D	A7	00	00	00	00	98	00	00	00	00	00	00	00
15531248	01	00	00	00	18	00	00	00	02	00	00	00	00	00	00	00
15531264	0F	00	0E	00	28	00	00	00	28	00	70	00	F8	10	01	00
15531280	00	00	88	02	00	00	08	00	00	00	00	00	00	00	00	00
15531296	AF	1C	00	00	00	00	00	00	64	32	02	00	00	00	02	00
15531312	70	00	5A	00	00	00	00	00	00	4C	01	00	00	00	02	00
15531328	32	87	FE	35	CC	26	D5	01	00	EC	13	A6	C4	0E	D5	01
15531344	52	87	EF	50	CE	26	D5	01	32	87	FE	35	CC	26	D5	01
15531360	00	00	01	00	00	00	00	00	32	F3	00	00	00	00	00	00
15531376	20	00	00	00	00	00	00	00	0C	02	42	00	52	00	41	00
15531392	49	00	4E	00	5F	00	7E	00	31	00	2E	00	4A	00	50	00
15531408	47	00	00	00	00	00	1D	00	B3	9F	1D	A7	00	00	00	00
15531424	9A	9F	1D	A7	00	00	00	00	9A	9F	1D	A7	00	00	00	00
15531440	A0	00	00	00	00	00	00	00	01	00	00	00	18	00	00	00
15531456	04	00	00	00	00	00	00	00	05	00	06	00	28	00	78	00

Fig:36 traces of transaction in Brain_scan.jpg

Workflow of deletion process has been illustrated in fig 38

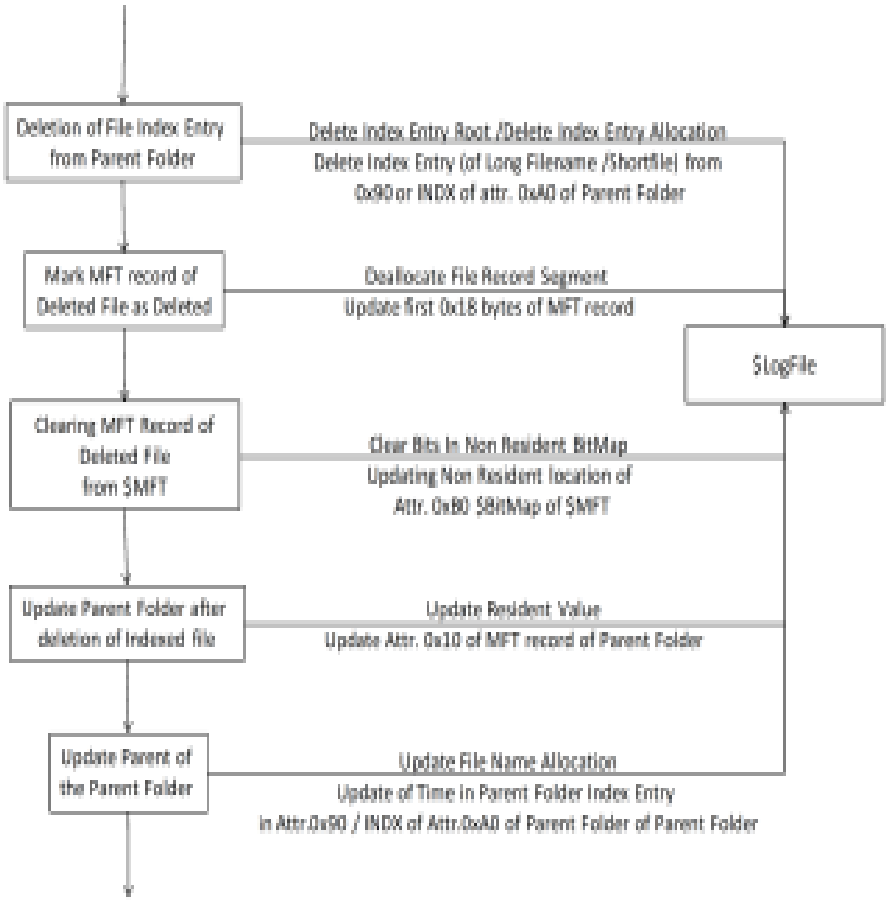


Fig 38: Deletion Process of a file in the $logfile

5.5 Summary-Forensic Artifacts

File creation and deletion process as discussed in chapter4 left many artifacts and they are backedup by Volume Shadow Copy.

Table 14: Forensic Artifacts of deleted files using Recycle.Bin

	$MFT		$Logfile	
Artifacts	**File**	**Parent**	**File**	**Parent**
File Name	Record of MFT remain unchanged Only the record header does not change Attr.0x30 (name started with $R followed by six random integers) Metadata of deleted MFT record is stored in $Ixxx file which containsthe name and location of the file.	Deleted folder and files are completely removed from theparent's folder and moved to the Recycle.bin folder of current user renaming to $Rxxx.Index entry will be renamed (name+time+size) of $Rxxx in Recycle bin of Current User Folder If original / renamed file was indexed at the end of 0x90 or INDX of Attr. 0xA0, deletion will leave this index entry (name+time+size) as file slack	$logfile records remain unchanged until overwritten by a new one when $logfile is out of space but during the deletionprocess, new records have been generated Delete Index Entry Root / Index Entry Allocation: Index Record of $Ixxx in Attr. 0x90/INDX of Attr. 0xA0 of Current User Folder of Recycle Bin	$logfile records remain unchanged until overwritten by a new one when $logfile is out of space but during the deletionprocess, new records have been generated Delete Index Entry Root / Index Entry Allocation: Index Record of $Rxxx in Attr. 0x90 / INDX of Attr.0xA0 of Current User Folder of Recycle Bin
File Time	Attr. 0x10 (Time Info related to entire MFT Record). This time will also give information on file deletion to Recycle Bin Attr. 0x30 (Time Info Related to File Name). This time will also give information of file deletion to Recycle Bin	Renamed Index Entry (name+time+size) of $Rxxx in Current User Folder of Recycle Bin If original / renamed file was indexed at the end of 0x90 or INDX of Attr. 0xA0, deletion will leave this index entry (name+time+size) as file slack	All $logfile records created during File Creation, remain unchanged. They are overwritten with new records after entire $Logfile space has been consumed. However, the deletion process generates its own records Delete Index Entry Root / Index Entry Allocation: Index Record of $Ixxx in Attr. 0x90/INDX of Attr. 0xA0 of Current User Folder of Recycle Bin	All $logfile records created during File Creation, remain unchanged. They are overwritten with new records after the entire$Logfile space has been consumed. However, the deletion process generates itsown recordsUpdate File Name Allocation: Time and Size info of Index entry of Current User Folder (longfilename and short film) in INDX of Attr. 0xA0 of Recycle Bin after clearing recycle bin reveals the time of recycle bin empty operation. Update Resident Value: Attr. 0x10 of parent folder after the deleted file has been de-indexedReveals the time of file deletion to recycle bin

During the experiment following important artifacts at following offsets found.

Table 15: Artifacts found in the experiment

Offset	Remarks	Findings/Artifacts
0x1E00	Start of the Shadow Service	Volume Shadow Service Header always starts from this offset
0x71f0000	Offset of the first catalog block offset	First catalog block
625f84000		The offset of first store block extracted from entry type of 0x03
625f8000	Relative offset of first store block	625f8C000
	Time Stamp Creation time of 1st Volume Shadow	*01D4BD41EBB41F96(5-2-2019 11:31)*
639f84000	Relative offset of second store block	639f84000+0xC000=639f84C00
	Creation time of 2nd Volume Shadow	*01D4BD42189D1FE6 (5-2-2019 11:43)*
64DF8400	Relative offset of third store block	64DF84000+0xC000 =64DF84000
	Creation time of 3rd Volume Shadow	*01D4BD424f738980(5-2-2019 11:57)*
661f84000	Relative offset of fourth store block	661f84000+0xc000=661f84c000
	Creation time of 4th Volume Shadow	Creation time of 4th Volume Shadow *01D4BD42943848F1(5-2-2019 12:16)*
C5FE5800	*Indexed attribute of deleted folder*	*S-1-5-21-3207469532-412700025-1003426619-1001~$I30 (90)*
C7A09920	*Offset of $I file of deleted folder Account_Info*	*$I8HLBDG file*
C5EAFDB8	*Account_information.txt File was permanently deleted from the system but present in the Volume Shadow Copy)*	*Account_information.txt file contents*
1FA237000	*Bank statement.png file which was deleted but recovered using Volume Shadow Copy)*	*Bank statement.png file*
0x64E75CDE0	*LSN of Account_info folder in $logfile*	*LSN of Account_info folder in $logfile*

After analyzing the Volume Shadow copy Recycle.Bin little difference has been found between $I file of $Recycle.Bin of Windows 7,8,10

Table 16: Comparison between $I file of Window 7,8 and 10

	Header (0x01)		Deleted File Size		Deleted Time Stamp		File Name Length		File Name Path	
	S.B*	Size	S.B*	Size	S.B*	Size	S.B*	Size	S.B*	Size
Windows 7,8	0	8	8	16	16	8	24			
Windows 10	0	8	8	16	16	8	24	4	2 8	Variable length

*Chapter*6

Review of tools used for the analysisof Volume Shadow Copy

After manual analysis of the Volume Shadow Copy for verification purpose, automated analysis of the Volume Shadow Copy has also been done. For automated analysis different freely available tools which support the Volume Shadow Copy have been used.

Many Digital Forensic Tools does not support Volume Shadow Copy.

Shadow Explorer

Shadow Explorer is a digital forensic tool specially made for analysis of Volume Shadow Copy. Allow the users of windows vista/7/8 and 10 to access the point in time shadow copies of the system.It does not support raw images. It only accesses the Volume Shadow Copies from the live volume. user can export any folder/file. [*23*] . The interface of the shadow explorer has been shown in Fig 37

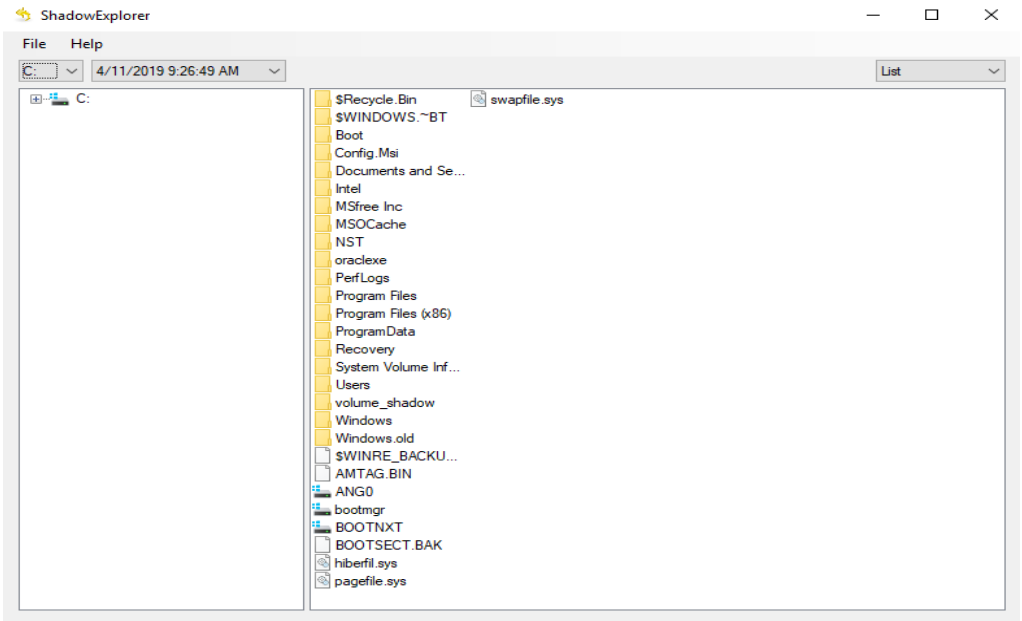


Fig: 37 Screen Shot of Shadow Explorer

Shadow Copy View

Shadow copy View is a simple tool that let the user explore the volume shadow copies of his system. It also does not support images of the system. It gives the edge to the investigator to see the directories before manual analysis of the volume Shadow copies [*26*]. The simple view of the Volume Shadow View is shown in Fig:33

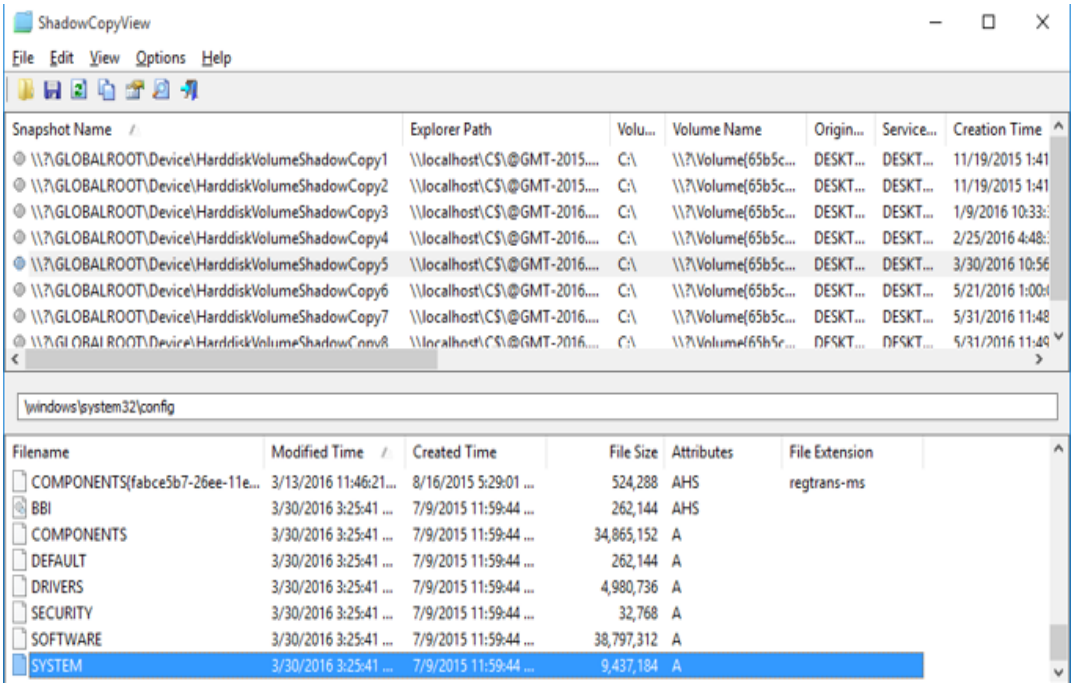


Fig 38: Screenshot of ShadowCopyView

Winhex

Winhex is a powerfuldigital forensic toolfor the investigators to investigate and find the evidence. Its academia version is freely available. In Winhex Refine Snapshot feature is included but only in licensed version under the specialist option shown in fig 39. It gives the option of a manual analysis of the Volume Shadow Copy briefly explained in the above chapters.

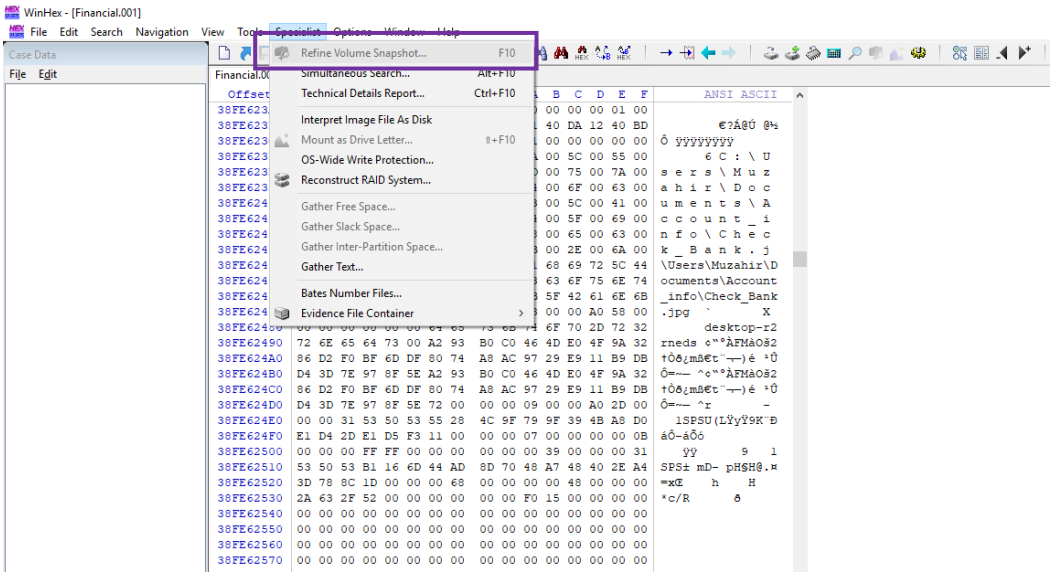


Fig 39: Screenshot of winhex academia version

Comparative analysis of these Forensic tools has been shown in table 13:

Table13: Comparative analysis of Volume Shadow Copy

Name	Freely available	Deals with Image files	System files	Shadow Copy Analysis facility
Shadow Explorer	Yes	No	Yes	No (Only gives the view of directory
ShadowCopyView	Yes	No	Yes	No (Only gives the view of directory)
Forensic Explorer	No (Only 30 days Trial with limited features)	Yes	Yes	Yes (Only in Licensed version)
WinHex	No (Only Academia Version is available)	Yes	Yes	Yes (Only in Licensed version)

C h a p t er 7

Discussion and Future work

This research thesis has focused on the Volume Shadow Copy its importance and way of analyzing the Volume Shadow Copy of windows 10. Volume Shadow Copy is a goldmine for forensic investigators as it contains an older version of the files with $logfile backed up with each Volume Shadow Copy and $Recycle.Bin both of these files are enriched with forensic artifacts.

It is essential for the computer forensic investigators to know about the importance and complications of the Windows Volume Shadow Copy to extract the evidence. Volume shadow copy contains the time in point snapshots of the system. Instead of saving the whole file at once it just saves the difference from the previously stored file. so, locating just only the difference file is a little bit tricky and complicated. It does not save the whole file every time the shadow is created. If any file is deleted and before deletion Volume Shadow Copy has been created it will be saved even if it is deleted bypassing the Recycle Bin. Volume Shadow Copy also contains the $logfile till that point of time provides the clumps of important data that can be investigated to find the evidence. Manual Analysis of the Volume shadow copy from a system image has been done which is accurate but lengthy and time-consuming process. Volume Shadow copy is a treasure for the forensic investigator, so more precise and accurate tools should be developed which can analyze the volume shadow copy accurately and precisely as most of the tools provide only the view of the volume shadow copies of live systems and do not provide the facility to analyze volume shadow copy from system images.

To conduct the experiment two case scenarios have been generated and tested. To successfully conduct the experiment different files have been created and modified and moved around different directories to create the traces. After taking the images they have been analyzed through Winhex Academia version.

Different shadow copies have been created for modified, changed the path and deleted files/folders. Different traces have been found in different shadow copies. The $logfile and $Recycle.Bin has also been analyzed. The changed state of the $logfile for that point of time has also been copied and analyzed. $Recycle.Bin is important as before emptying the recycle.bin if it is copied by the Volume Shadow Service, it's all contents will also be saved and can be a great point of interest for the forensic investigators.

In this thesis, only $logfile and $Recycle.bin has been analyzed but Volume Shadow Copy has a lot more than these two files. It not only has the previous versions of the user files but also has the previous versions of the system files. Complete analysis of Volume Shadow Copy with all feature e.g $logfile in each shadow can be taken as a continuation of this research work.

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